



Open Metering System Specification

Security Mode 13 - Transport Layer Security (TLS)

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F.1 Normative references

- [OMS-S2] OMS-Specification Volume 2 Issue 4.2.1 (the superior document)
- [BSITR03109-1] BSI Technische Richtlinie BSI TR-03109-1: Anforderungen an die Interoperabilität der Kommunikationseinheit eines intelligenten Messsystems] Version 1.0.1, Datum 16.01.2019
<https://www.bsi.bund.de/SharedDocs/Downloads/DE/BSI/Publikationen/TechnischeRichtlinien/TR03109/TR03109-1.pdf>
- [BSITR03116-3] BSI Technische Richtlinie TR-03116-3 Kryptographische Vorgaben für Projekte der Bundesregierung; Teil 3: Intelligente Messsysteme, Datum 23.04.2018
<https://www.bsi.bund.de/SharedDocs/Downloads/DE/BSI/Publikationen/TechnischeRichtlinien/TR03116/BSI-TR-03116-3.pdf>
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- [EN13757-4:2019] Communication systems for meters — Part 4: Wireless M-Bus communication;
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The AES-CMAC Algorithm; June 2006
<https://tools.ietf.org/html/rfc4493>
- [RFC5246] Network Working Group Request for Comments: 5246
The Transport Layer Security (TLS) Protocol Version 1.2; August 2008
<https://tools.ietf.org/html/rfc5246>
- [RFC5289] Network Working Group Request for Comments: 5289
TLS Elliptic Curve Cipher Suites with SHA-256/384 and AES Galois Counter Mode (GCM); August 2008
<https://tools.ietf.org/html/rfc5289>
- [RFC6066] Internet Engineering Task Force (IETF) Request for Comments: 6066
Transport Layer Security (TLS) Extensions: Extension Definitions; January 2011
<https://tools.ietf.org/html/rfc6066>
- [RFC7027] Internet Engineering Task Force (IETF) Request for Comments: 7027
Elliptic Curve Cryptography (ECC) Brainpool Curves for Transport Layer Security (TLS); October 2013
<https://tools.ietf.org/html/rfc7027>
- [RFC7366] Internet Engineering Task Force (IETF) Request for Comments: 7366
Encrypt-then-MAC for Transport Layer Security (TLS) and DTLS, September 2014
<https://tools.ietf.org/html/rfc7366>
- [RFC8422] Network Working Group Request for Comments: 8422
Elliptic Curve Cryptography (ECC) Cipher Suites for Transport Layer Security (TLS) Versions 1.2 and Earlier; Aug. 2018
<https://tools.ietf.org/html/rfc8422>

F.2 Introduction

This annex is an amendment to the OMS-Specification [OMS-S2] and specifies the security mechanisms of security mode 13 for the M-Bus protocol.

Security mode 13 requires asymmetric encryption and specifies the utilization of Transport Layer Security (TLS). Since TLS is a security implementation with many configuration possibilities, this specification provides limitations, extensions and options to TLS. This includes also an applicable selection of crypto and cipher suites.

The management of security related services and the exchange of security information is handled with the SITP protocol [EN13757-7:2018]. E.g., the SITP protocol provides procedures for key and certificate updates and exchanges.

This specification comes with examples for the frame structure when using the AFL in combination with TLS, key exchanges and TLS messages. The examples are available for M-Bus and wM-Bus. Further examples for TLS channel establishment management, application data transfer and master key updates are defined utilizing the SITP protocol. Also, an example of a certificate is present.

Following special terms are used in this annex:

TLS session: Security association between two TLS endpoints. The TLS session starts with negotiation of secrets between the TLS endpoints during a non-resumed TLS handshake. The TLS session ends either by session timeout or by replacement with a new session between the same endpoints.

TLS channel: Reliable transport established between two TLS endpoints. The TLS channel is established with a response to a TLS ChannelRequest message. The TLS channel is closed either by end of the session or by transmission of a TLS alert. A TLS channel uses the AFL.

TLS connection: This term is not used in this document, because neither DLL/ELL nor AFL is connection oriented.

TLS record Is either a TLSPlaintext or TLSCiphertext.

TLSPlaintext Contains a TLS record header and content specific data.

TLSCiphertext Contains a TLS record header and TLS cipher suite specific header, content specific data and TLS cipher suite specific trailer.

The following table list the terms used in this specification and its corresponding name used in [RFC5246].

Table F.1 – List of used in [RFC5246]

Name in this Annex	Name in [RFC5246]
ContentType	TLSPplaintext.type TLSCiphertext.type
ProtoMajor	TLSPplaintext.version.major TLSCiphertext.version.major
ProtoMinor	TLSPplaintext.version.minor TLSCiphertext.version.minor
TLSSDULen ^a	TLSPplaintext.length TLSCiphertext.length
HSMsgType	Handshake.msg_type
HSMsgLength	Handshake.length
HSClientVersionMajor	Handshake.ClientHello.client_version.major
HSClientVersionMinor	Handshake.ClientHello.client_version.minor
HSServerVersionMajor	Handshake.ServerHello.server_version.major
HSServerVersionMinor	Handshake.ServerHello.server_version.minor
HSClientRandom	Handshake.ClientHello.random
HSServerRandom	Handshake.ServerHello.random
HSSessionIDLength	length(Handshake.ClientHello.session_id)
^a Length of TLS service data unit	

F.3 Asymmetric encryption with security mode 13

F.3.1 Required TLS operation for security profile C

Asymmetric encryption is performed with security profile C (security mode 13) which allows the usage of TLS for the M-Bus protocol.

5 To be able to implement the TLS protocol as defined in [RFC5246] in an environment with such limited networking resources as wireless M-Bus, the following limitations and extensions are required.

- OMS end-devices shall implement the TLS client functionality; gateways shall implement the TLS server functionality.

10 • TLS1.2 shall be supported by client and server. In the future also higher TLS versions should be supported.

- Servers shall support the ClientHello MaxFragmentLength extension (see [RFC6066]) with min. 512 bytes fragment size.

NOTE 1: This has not to be confused with the AFL-fragment length.

15 **NOTE 2:** In future version of this Annex it is intended to improve fragmentation according to RFC8449.

- In addition to the standard HMAC size of 32 bytes (as defined in [RFC5246]), servers shall support the ClientHello TruncatedHmac extension (see [RFC6066]) to send the first 10 bytes of the negotiated HMAC and verify the first 10 bytes of the calculated hash.
- 20 Clients may support the ClientHello TruncatedHmac extension (see [RFC6066]) in order to save data overhead.

- Servers shall accept the EncryptThenMac ClientHello extension and support EncryptThenMac according to [RFC7366] if AES-CBC cipher suites are offered

25 • Clients should send the EncryptThenMac ClientHello extension and should support EncryptThenMac according to [RFC7366] if an AES-CBC cipher suite is requested

- Clients shall accept the additional TLS ChannelRequest Message. This message shall be authenticated with at least 8 bytes of AES128 CMAC protection.

- Handshake messages (according to Table F.12) shall be transmitted/ accepted by the client (OMS end-device) only in the required sequence of RFC5246, Figure1 and Figure 2.

30 • TLS session renegotiation is not allowed.

- A TLS session can be resumed. This document allows a max. TLS session lifetime of 1 month (31 days) unless not more than 5 Mbytes of TLS data (according to [BSI-TR03116-3]) is transmitted in both directions within this TLS session. TLS session resumption shall be implemented in the gateway.

35 • A new (authenticated) TLS ChannelRequest from server (gateway) to client (OMS end-device) shall close any existing TLS channel between client (OMS end-device) and server (gateway) and establish (resume or negotiate) a new one.

- For TLS crypto suites based on elliptical curves, the uncompressed point format shall be used.

40 • Clients and servers shall support ECDSA for signature validation and generation with the required hash algorithms according to [BSI TR-3116-3], 6.4.1.

- Clients and servers shall support the `supported_signatures_extension` according to [RFC5246] 7.4.1.4.1 to indicate the supported signature algorithms (e.g. ECDSA) and hash algorithms.

45 TLS crypto suites based on elliptical curves shall be implemented according to [RFC8422] or [RFC5289]

The identifiers for brainpool TLS named curve IDs are defined in [RFC7027]. At least the brainpoolP256r1 and Secp256r1/NIST P256 shall be supported for ECDHE and ECDSA keys. The support of the other TLS named curves is recommended.

NOTE 1: The default FragmentSize limit in TLS is 16 kbytes (which shall be buffered by the receiver and the transmitter). For memory size restricted OMS end-devices a max_fragment_length extension can be negotiated, in case the client also supports it. The negotiation is defined in [RFC6066].

NOTE 2: TLS1.2 suggests but does not require a limitation of the (resumeable) TLS session lifetime. For security reasons (for example the derived ephemeral keying material may be compromised by memory reads) the TLS session life time shall be limited.

NOTE 3: A session renegotiation is not necessary and shall not be allowed. With nearly the same number of handshake messages and data, the existing session can be closed and a new session can be established.

NOTE 4: After termination of the session (in an orderly manner the gateway shall send a CloseNotify message to the OMS end-device), the OMS end-device may immediately start a new TLS session by sending a ClientHello to the gateway.

Table F.2 – List of supported elliptical curves

Curve name	TLS named curve ID
brainpoolP256r1	001Ah
brainpoolP384r1	001Bh
brainpoolP512r1	001Ch
Secp256r1/NIST P256	0017h
Secp384r1/NIST P384	0018h

NOTE: The [BSITR03116-3] recommends several cipher suites and its usage period. It is recommended to consider this technical guideline for the selection of supported cipher suites. This technical report will be updated periodically based on the [BSITR03116-3]!

F.3.2 Transport layer / configuration field

The TLS security mode is declared with a 13 (0Dh) in the field “security mode” (MMMMM).

NOTE: Bits for link control are served by the Extended Link Layer (ELL), which has always been applied together with TLS.

The configuration field (CF) of the security mode 13 is defined in Table F.3:

Table F.3 – Configuration field of security mode 13

Bit 15	Bit 14	Bit 13	Bit 12	Bit 11	Bit 10	Bit 9	Bit 8	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0
Content of Message	Content of Message	Reserved	Security Mode	Security Mode	Security Mode	Security Mode	Security Mode	Number of bytes for Mode 13	Number of bytes for Mode 13	Number of bytes for Mode 13	Number of bytes for Mode 13	Number of bytes for Mode 13	Number of bytes for Mode 13	Number of bytes for Mode 13	Number of bytes for Mode 13
C	C	0	M	M	M	M	M	N	N	N	N	N	N	N	N

M shall always be 13 (0Dh) to declare an encryption with TLS

NOTE: The applied TLS version has to be traced from TLS header
C declares the content of the message and shall comply with [EN13757-3:2018].

N The N field is not used in security mode 13 and shall always be set to FFh. The number of following bytes is taken from field TLSSDULen in the TLS record header (see Table F.9).

F.3.3 Transport layer / configuration field extension

- 5 Security mode 13 requires a Configuration Field Extension (CFE) with a size of one byte according to Table F.4.

Table F.4 – Configuration Field Extension of security mode 13

Bit 23	Bit 22	Bit 21	Bit 20	Bit 19	Bit 18	Bit 17	Bit 16
Reserved	Reserved	Reserved	Reserved	ProtocolType 3	ProtocolType 2	ProtocolType 1	ProtocolType 0
0	0	0	0	P	P	P	P

P is the Protocol Type as defined in Table F.5

Table F.5 – Protocol types of security mode 13

Content of protocol type P	Meaning
0000	TLSSCHAN (TLS ChannelRequest)
0001	TLSPROT (TLS handshake , TLS alert, TLS application (see Table F.6))
0010..1111	Reserved

10 F.3.4 Structure of different TLS messages

F.3.4.1 General

This section presents the structure and the applicable content types of TLS. The Table F.6 shows the different content types of the TLS protocol. The following sections show the structure of the message depending on the content type.

15 **Table F.6 – Mapping between TLS protocol type and ContentType**

Symbol name	ContentType name	ContentType (according [RFC 5246])	TLS protocol type (according to Table F.5)
TLSSCHAN	TLS ChannelRequest	00 _h	TLSSCHAN
TLSSHS	TLS Handshake	14 _h , 16 _h	TLSPROT
TLSSALERT	TLS Alert	15 _h	TLSPROT
TLSSAPP	TLS Application	17 _h	TLSPROT

The TLS structure follows the CF and CFE bytes in case of security mode 13. It corresponds with the “optional TPL header fields” and (depending on the ContentType) with “Encrypted/unencrypted application data” and “TPL trailer field” as shown in [EN13757-7:2018] Table 20.

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Table F.7 – General structure of a TLS Message

		TLCHAN	TLHS, TLALERT c	TLHS, TLALERT d	TLAPP
TPL-Header	CI-field	5F _h	5F _h , 9E _h , 9F _h	5F _h , 9E _h , 9F _h	e.g. 5B _h , 7A _h , C3 _h , C4 _h
	CI-dependent fields	See EN 13757-7 Table 11 + 12	See EN 13757-7 Table 11 + 12	See EN 13757-7 Table 11 + 12	See EN 13757-7 Table 11 + 12
	Configuration field extension	See Table F.4	See Table F.4	See Table F.4	See Table F.4
	Optional TPL header fields ^a	TLCHAN header (see Table F.8)	TLS record header ^b (see Table F.9)	TLS record header ^b (see Table F.9)	TLS record header (see Table F.9)
		No TLS cipher suite specific header	No TLS cipher suite specific header ^b	TLS cipher suite specific header ^b	TLS cipher suite specific header
APL data	optional	No data	Content specific data ^b	Content specific data ^b	APL Data
TPL-Trailer	optional	No TLS cipher suite specific trailer	No TLS cipher suite specific trailer ^b	TLS cipher suite specific trailer ^b	TLS cipher suite specific trailer
^a The Optional TPL header fields need to be present for all message types used in this annex! ^b To minimize data overhead several TLS records can be concatenated into one TLS message. ^c TLSPplaintext (see F.2) ^d TLSCiphertext (see F.2)					

NOTE: This presentation only shows transport and application layer

Table F.8 – TLSCHAN header

TLS record header of TLSCHAN			
ContentType (1 byte)	reserved ^b (1 byte)	reserved ^b (1 byte)	TLSSDULen ^a (2 byte)
^a TLSSDULen is transmitted with MSB first			
^b TLSCHAN don't use fields ProtoMajor and ProtoMinor			

Table F.9 – TLS record header according to TLS1.2 [RFC 5246]

TLS record header of TLSALERT, TLSHS, TLSAPP			
ContentType (1 byte)	ProtoMajor (1 byte)	ProtoMinor (1 byte)	TLSSDULen ^a (2 byte)
^a TLSSDULen is transmitted with MSB first			

ContentType shall be according to Table F.6.

ProtoMajor and ProtoMinor shall be according to [RFC 5246] (see Table F.1)

5 TLSSDULen is the length (in bytes) of the TLSPlaintext.length or TLSCiphertext.length and shall be used according to [RFC 5246] (see also Table F.1)

F.3.4.2 TLS ChannelRequest

Table F.10 – TLS ChannelRequest

CI = 5Fh	ADDR, ACC, ST, CF = 13; TLSCHAN	TLS ContentType (00h)	Reserved (00h)	Reserved (00h)	TLSSDULen MSB = 00h	TLSSDULen LSB = 00h
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This content type requires always a CMAC protection by the AFL layer.

F.3.4.3 TLS Handshake

Table F.11 – TLS Handshake

CI = 5Fh/ 9Eh/ 9Fh	[ADDR,]ACC, ST, CF = 13; TLSPROT	TLS ContentType (14h, 16h)	TLS ProtoMajor (03h)	TLS ProtoMinor (03h)	TLSSDULen MSB	TLSSDULen LSB
HSMType	TLS MsgData					

The following table describes the used handshake message types according to RFC5246 and the support of CMAC-protection.

Table F.12 – Handshake message types

ContentType	HSMType	Name of HSMType in RFC	CMAC required
14h	01h	ChangeCipherSpec	NO
16h	01h	ClientHello	YES
16h	02h	ServerHello	NO
16h	0Bh	Certificate	NO
16h	0Ch	ServerKeyExchange	NO
16h	0Dh	CertificateRequest	NO
16h	0Eh	ServerHelloDone	NO
16h	0Fh	CertificateVerify	NO
16h	10h	ClientKeyExchange	NO
16h	14h	Finished	NO

F.3.4.4 TLS Alert

Table F.13 – TLS Alert with HMAC

CI = 5Fh/ 9Eh/ 9Fh.	[ADDR,]ACC, ST,CF = 13; TLSPROT	TLS ContentType (15h)	TLS ProtoMajor (03h)	TLS ProtoMinor (03h)	TLSSDULen MSB	TLSSDULen LSB
TLS random IV (16 bytes)	Encrypted alert data ^a				Padding ^b	TLS HMAC (32 bytes)
^a The encrypted alert data contains a 2 byte content according to chapter 7.2 of [RFC5246].						
^b Padding according to [RFC5246]						

Table F.14 – TLS Alert with GCM

CI = 5Fh/ 9Eh/ 9Fh.	[ADDR,]ACC, ST, CF = 13; TLSPROT	TLS ContentType (15h)	TLS ProtoMajor (03h)	TLS ProtoMinor (03h)	TLSSDULen MSB	TLSSDULen LSB
GCM nonce	Encrypted alert data ^{a, b}					
^a The encrypted alert data contains a 2 byte content according to chapter 7.2 of [RFC5246].						
^b The padding bytes and the authentication tag are part of the encrypted alert data.						

During a valid TLS session, the TLS Alert is protected by TLS. In all other cases the TLS Alert shall be protected by the CMAC in the AFL layer.

F.3.4.5 TLS Application

Table F.15 – TLS Application with HMAC

CI = 5Fh/ 9Eh/ 9Fh.	[ADDR,]ACC, ST, CF = 13; TLSPROT	TLS ContentType (15h)	TLS ProtoMajor (03h)	TLS ProtoMinor (03h)	TLSSDULen MSB	TLSSDULen LSB
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CI = 5Bh/ 7Ah/ 72h AppData	[ADDR,]ACC, ST, CF = 13; TLSPROT	TLS ContentType (17h)	TLS ProtoMajor (03h)	TLS ProtoMinor (03h)	TLSSDULen MSB	TLSSDULen LSB
TLS random IV (16 bytes)	Encrypted app. data TLS HMAC (32 bytes)					TLS HMAC (32 bytes)

This content type is protected by TLS and shall not apply a CMAC protection.

Table F.16 – TLS Application with GCM

CI = 5Bh/ 7Ah/ 72h AppData	[ADDR,]ACC, ST, CF = 13; TLSPROT	TLS ContentType (17h)	TLS ProtoMajor (03h)	TLS ProtoMinor (03h)	TLSSDULen MSB	TLSSDULen LSB
GCM nonce	Encrypted app. data					
^a The encrypted app. data contains the padding bytes and the authentication tag beside the application data.						

NOTE: The maximum size for an AFL fragmented message is limited to 16 kbytes. Therefore the maximum TLSSDULen is less than 16kbytes, depending on the other used fields, e.g. in the TLS header. Client and server can agree to use a smaller maximum TLS fragment size (ClientHello extension).

Be aware that existing TLS 1.2 [RFC5246] with AES-CBC cipher suite applies HMAC over the plain message. The encryption is done thereafter. OMS end-device and gateway are required to support the encrypt-then-MAC according to [RFC7366] for AES-CBC cipher suites.

F.3.5 TLS message sequence

F.3.5.1 Establishment of TLS channel and prevention of ClientHello flooding

Most wireless M-Bus channels are limited in bandwidth and in duty cycle. Sending many ClientHello will rapidly reduce the transmit credits of the sender or fill the physical medium. This is true for both the master and the slave M-Bus device (gateway and OMS end-device). The sender may limit its transmission rate to prevent a DoS attack to the bandwidth limited medium.

Another problem is the (battery-) power consumed by calculation of the random number, triggered by unnecessary ClientHellos.

To take measures against these problems the TLS role “client” is assigned to the OMS end-device and the TLS role “server” is assigned to the gateway.

The OMS end-device shall only accept TLS ChannelRequest from the gateway if the TLS ChannelRequest received is AFL CMAC authenticated and verified.

The TLS channel starts, when the client (OMS end-device) sends a TLS ClientHello handshake message to the server (gateway) (TLS 1st MsgBlock). The server (gateway) shall only accept ClientHello from the client (OMS end-device) if the received ClientHello is AFL CMAC authenticated and verified.

The server (gateway) answers with ServerHello and ServerHelloDone within a TLS message with typically more than 400 bytes length (TLS session not resumed). According to RFC[5246], the server (gateway) shall answer with ServerCertificate, ServerKeyExchange and ClientCertificateRequest, if the related requirements of the TLS configuration are fulfilled.

Any existing TLS channel between server (gateway) and client (OMS end-device) shall be silently closed in the client (OMS end-device), after receiving an authenticated and verified TLS ChannelRequest from the server (gateway). Any existing TLS channel between server and client shall be silently closed in the server, before sending an authenticated TLS ChannelRequest to the client. A TLS channel shall be identified by the Application Layer Address ADDR of the server (consists of ID, manufacturer, version, device type).

NOTE: The following message flights uses a SND-UD for data transfer to the OMS end-device. The gateway may also apply a SND-UD2 for non-fragmented wireless M-Bus messages (see [OMS-S2], 5.2.3).

F.3.5.2 TLS session with full handshake, initiated by TLS server (gateway)

Table F.17 – TLS session with full handshake, initiated by TLS server (gateway)

Gateway (LLA = COM)	Message flight	OMS end-device (LLA = MTR)
TLS ChannelRequest to enter frequent access cycle with new TLS channel	SND-UD(C = 73h); CI = 8Ch; CI = 90h; CMAC; CI = 5Fh; MTR; CF(MODE 13; TLSCHAN) ACK(C = 00h); CI = 8Ch; CI = 8Ah 	The client enters frequent access cycle and requests a new TLS-channel. Waits for REQ-UD2 from server to send ClientHello
The server requests to receive	REQ-UD2(C = 5Bh); CI = 8Ch; CI = 80h; MTR RSP-UD(C = 08h); CI = 8Ch; CI = 90h; CMAC; CI = 9Eh; Flight 1	Client sends ClientHello (without SessionID)
Server sends ServerHello, Certificate, ServerKeyExchange, CertificateRequest and ServerHelloDone in fragmented messages.	SND-UD(C = 73h); CI = 8Eh; CI = 90h; FragID = 1; TLS RecLen; CI = 5Fh; CF(MODE 13; TLSHS) ACK(C = 00h); CI = 8Eh SND-UD(C = 53h); CI = 8Eh; CI = 90h; FragID = 2; ACK (C = 08h); CI = 8Eh; CI = 8Ah Flight 2 ^a	
Server requests to receive	REQ-UD2(C = 7Bh); CI = 8Ch; CI = 80h RSP-UD(C = 08h); CI = 8Eh; CI = 90h; FragID = 1; TLS RecLen; CI = 9Eh; CF(MODE 13; TLSHS) REQ-UD2(C = 5Bh); CI = 8Eh RSP-UD(C = 08h); CI = 8Eh; CI = 90h; FragID = 2 Flight 3 ^a	Client sends Certificate, ClientKeyExchange, CertificateVerify, ChangeCipherSpec, Finished
Server sends ChangeCipherSpec, Finished.	SND-UD(C = 73h) ; CI = 8Ch ; CF(MODE 13; TLSHS) ACK(C = 00h); CI = 8Ch; CI = 8Ah Flight 4	

^a Due to the length of the certificate it might be necessary to use a third fragment.

All numbers are hexadecimal.

F.3.5.3 Resumed TLS session, initiated by TLS server (gateway)

Table F.18 – Resumed TLS session, initiated by TLS server (gateway)

Gateway (LLA = COM)	Message flight	OMS end-device (LLA = MTR)
Optional immediate SND-UD (for OMS end-devices with receiver always on) The gateway sends TLS ChannelRequest	SND-UD(C = 53h); CI = 8Ch; CI = 90h; CMAC; CI = 5Fh; CF(MODE 13; TLSCHAN) →	Since the receiver is not always open, the message is not received.
	SND-NR(C = 44h); CI = 7Ah ←	OMS end-device sends periodic, encrypted data to gateway. Payload AES encrypted and opens receiver after 2 to 3 ms.
SND-UD to enter the frequent access cycle with a TLS ChannelRequest	SND-UD(C = 73h); CI = 8Ch; CI = 90h; CI = 5Fh; CF(MODE 13; TLSCHAN) → ACK(C = 00h); CI = 8Ah ←	The OMS end-device enters frequent access cycle and requests a new TLS channel. Waits for REQ-UD2 from gateway to send ClientHello
	REQ-UD2(C = 5Bh); CI = 8Ch; CI = 80h → RSP-UD(C = 08h); CI = 8Ch; CI = 8Eh CF(MODE 13; TLSHS) ← Flight 1	OMS end-device sends ClientHello with SessionID
Gateway sends a ServerHello with SessionID (reflects the SessionID received in the ClientHello message) and ChangeCipherSpec, finished in one message.	SND-UD(C = 73h); CI = 8Ch; CI = 5Fh; CF(MODE 13; TLSHS) → ACK(C = 00h); CI = 8Ah ← Flight 2	
Gateway requests to receive	REQ-UD2(C = 5Bh); CI = 8Ch; CI = 80h → RSP-UD(C = 08h); CI = 8Ch; CI = 9Eh; CF(MODE 13; TLSHS) ← Flight 3	OMS end-device sends ChangeCipherSpec, Finished

All numbers are hexadecimal.

In case of TLS session resumption, the server shall ignore the supported elliptic curves extension and the supported point formats extension appearing in the current ClientHello message.

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F.3.5.4 Exchange of application data

Table F.19 – Exchange of application data

Gateway (LLA = COM)	Message flight	OMS end-device (LLA = MTR)
TLS encapsulated request	<p>SND-UD(C = 73h); CI = 8Ch; CI = 5Bh; CF(MODE 13; TLSAPP)</p> <p>→</p> <p>ACK(C=00h); CI=8Ch; CI=8Ah; CF=0</p> <p>←</p> <p>Flight x</p>	
	<p>or</p> <p>REQ-UD2(C = 53h); CI = 8Ch; CI = 80h</p> <p>→</p> <p>RSP-UD(C = 08h); CI = 8Ch; CI = 72h CF(MODE 13; TLSAPP)</p> <p>←</p> <p>Flight y</p>	Send TLS encapsulated response

All numbers are hexadecimal.

F.3.5.5 Terminate TLS channel

- 5 Because the TLS channel may be silently terminated (closed) by either client or server (for example with a reboot and power-loss) the client shall accept an (authenticated) TLS ChannelRequest to close an existing TLS channel to the sender of the TLS ChannelRequest and initiate a new TLS channel by sending a ClientHello to the sender of the TLS ChannelRequest.
- 10 The wireless M-Bus does not define a connection-oriented layer. Therefore, termination of the TLS channel cannot close the underlying connection layer. TLS channel terminated by server:

Table F.20 – TLS channel terminated by server

Gateway (LLA = COM)	Message flight	OMS end-device (LLA = MTR)
Sends TLS CloseNotify	<p>SND-UD(C = 73h); CI = 8Ch; CI=5Fh; CF(MODE 13; TLSALERT)</p> <p>→</p> <p>ACK(C = 00h); CI = 8Ch; CI = 8Ah</p> <p>←</p> <p>Flight <i>n-3</i></p>	
Requests close notify from client	<p>REQ-UD2(C = 5Bh); CI = 8Ch; CI = 80h</p> <p>→</p> <p>RSP-UD(C = 08h); CI = 8Ch; CI = 9Eh; CF(MODE 13; TLSALERT)</p> <p>←</p> <p>Flight <i>n-2</i></p>	Send TLS CloseNotify
At the end of communication, the DLL may terminate the Frequent Access Cycle	<p>SND-NKE(C = 40h); CI = 8Ch; CI=80h</p> <p>→</p> <p>Flight <i>n-1</i></p>	

All numbers are hexadecimal.

F.3.6 OMS end-device behaviour for different TLS channel states

F.3.6.1 OMS end-device messages outside the TLS channel

The OMS end-device needs to transmit periodically messages to allow an access to the OMS end-device (see [OMS-S2], 4.3.3). This can be message type SND-NR or ACC-NR (see [OMS-S2], 5.2.3). Both message types are transmitted unsolicited with no acknowledgement and may be lost. For that reason, the unsolicited OMS end-device message shall not be transmitted with security profile C (not within the TLS channel).

F.3.6.2 TLS channel usage for OMS end-devices conforming to Annex E.1

If an OMS end-device or gateway complies with [OMS-S2], Annex E clause E.1 it need to transfer all application data via the TLS channel (security profile C, see [OMS-S2], 9.1).

Under certain conditions a gateway may also accept unidirectional messages from OMS end-devices with security profile B (see “LKS2” in [BSITR03109-1]) when the TLS channel is not available. For that reason, the OMS end-device shall transmit a SND-NR or ACC-NR periodically in synchronised transmission slots (according to [OMS-S2], 4.3.2.1). The SND-NR shall contain simple consumption data (at least the current consumption values). It shall be protected with security profile B (according to [OMS-S2], 9.1). The alternative ACC-NR (see [OMS-S2], 5.2.3) contains no application data and applies no security profile.

NOTE 1: If the OMS end-device applies only the ACC-NR the gateway will never receive consumption data in case of disturbed TLS channel.

All other services like selection and request of data or the acceptance of commands shall be rejected in the case the message was not protected with security profile C (TLS). The OMS end-device shall respond an application error in the concerning response (see F.3.6.4).

NOTE 2: This rule requires that a client supports a key pair and a valid certificate by default to enable a bidirectional data transfer with the server. If the security information is not available, they need to be provided using the local service interface of the OMS end-device (see F.4).

F.3.6.3 TLS channel usage for OMS end-devices not conforming to Annex E.1

For OMS end-devices which not need to comply with Annex E, E.1 the use of a TLS channel (security profile C, see [OMS-S2], 9.1) is optional.

F.3.6.4 Behaviour in case of application error

If the client receives a command via the TLS channel and an error has happened, it has to respond an application error. Depending on the error condition, the error code is transmitted either encrypted (within the TLS channel) or plain (outside the TLS channel). The conditions are defined in Table F.21.

Table F.21 – Encryption and authentication of application errors

Error condition	Application error	Encryption of application error	Authentication of application error
Application command has been successfully received via TLS channel but not executed (e.g. command unknown or missing parameters)	Application error according to [EN13757-3:2018], Tab.21	TLS ^a	TLS ^a
Application command has not been successfully received, because decryption or authentication has failed (e.g. wrong key, invalid MessageCounter or TLS handshake has not been finished)	Application error 20h	No ^c	No ^c
Application command has not been successfully received, because the selected security mode is not supported (e.g. wrong security profile)	Application error 21h	No	AFL CMAC ^b
Application command has been successfully received but it was rejected, because it has been received outside the TLS channel.	Application error 22h	No	AFL CMAC ^b
^a Security profile C ^b Security profile B ^c No Security profile			

An application error shall not cause a finish of an established TLS session.

NOTE 1: The application errors 20h to 22h describes failure in the transport layer in reason of incorrect keying material or wrong security method. For that reason the security method is omitted when transmitting the command's denial.

NOTE 2: [OMS-S2], 4.3.4 describes how an OMS end-device protects against unauthorized communication requests.

NOTE 3: A gateway according [BSITR03109-1] does not accept any message without authentication being checked. The application error may be analysed by a radio logger.

F.4 Update of security information

F.4.1 General

This section describes how to update the security information in the OMS end-device and the gateway. It describes how to read and write new security information or how to trigger the generation of new security information in the OMS end-device.

The reading of security information from the OMS end-device is limited. Private keys and symmetric keys shall not be readable over any external interface of the OMS end-device.

The SITP allows transmitting of several blocks within one message. For this Annex the execution shall be applied as a combined transaction (either for all blocks or for none). This means the receiver needs to check all blocks first before executing the first block. If an error is detected in any block then the transaction (of all blocks) shall be rejected.

F.4.2 Master key update by the gateway

The master key update shall be performed using the SITP protocol (see Appendix F.A) over TLS (security mode 13). The new master key will be calculated by the OMS end-device with a random z_1 (a 128 bit random number) which has to be provided by the gateway.

NOTE: The installation of an already used MK shall be avoided.

F.4.2.1 SITP parameters during Master key update

KeyID

According to [OMS-S2], 9.2.3.2 the master key applies KeyID = 00h.

KeyVersion

The Initial KeyVersion of the Master key MK_0 after Reset to MK_0 is 00h. For each master key update the new KeyVersion shall be different to the currently active KeyVersion. It is recommended to increment it by 1. The KeyVersion FF_n is used as a wildcard and shall not be assigned as a discrete value.

In the following section MK is the current master key (version is n), MK' is the new master key (version after increment is n').

1. Before transferring the random z_1 , the gateway initiates a TLS channel between OMS end-device (TLS client) and gateway (TLS server). The AFL CMAC is required for TLS ChannelRequest and TLS ClientHello is computed using the current MK.
2. The following SITP commands are sent within TLS application records (security mode = 13h)
3. The gateway generates a new random z_1 for the OMS end-device and stores MK' without overwriting MK. In case of a reboot of the gateway between step 7 and 10, the MK' is required.

4. Transfer of the 16 byte random z1 into the OMS end-device is done using SITP with the following parameters:

Table F.22 – SITP parameters for master key update, step 4

Parameter	Value	Description
BCF	00 _h	Command “Transfer security information”
DSI	01 _h	Wrapped data structure according to NIST SP 800-38F type KWP
DSH1	FF _h	KeyID not used (no wrapping applied)
DSH2	FF _h	KeyVersion not used (no wrapping applied)
Key	z1	128 bit random z1
TargetTime	3080000000 _h	set to “invalid” as separate activation service is used
KeyID	00 _h	master key
KeyVersion	n’ ^a	according MK’
^a In case, the OMS end-device receives a KeyVersion = FF _h (unused) the OMS end-device computes the new KeyVersion n’ on its own by incrementing n.		

5. In case, the OMS end-device receives a KeyVersion equal to its current KeyVersion (n’ = n), an SITP status response is sent by the OMS end-device and the KeyUpdate is aborted with a permanent failure. The current master key MK is still in use.

Table F.23 – SITP parameters for master key update, step 5

Parameter	Value	Description
BCF	80 _h	Response “Transfer security information”
DSI	22 _h	Status response structure
DSH1	FF _h	KeyID not used
DSH2	FF _h	KeyVersion not used
Status response byte	21 _h	“Unknown or invalid KeyID/KeyVersion”

6. In case, the gateway does not receive a successful SITP status response, the KeyUpdate is aborted with a permanent failure. The current master key MK is still in use.

Table F.24 – SITP parameters for master key update, step 6

Parameter	Value	Description
BCF	80 _h	Response “Transfer security information”
DSI	22 _h	Status Response structure
DSH1	FF _h	KeyID not used
DSH2	FF _h	KeyVersion not used
Status response byte	00 _h	“Successful SITP command”

7. If the transfer has been successful, the OMS end-device calculates MK’ = CMAC(MK, z1) according to TR-03116-3, 7.1.1 and stores the new master key MK’ under KeyVersion n’, but does not activate the key.

8. Activation of the new master key and deactivation of the old master key is done using SITP with the following parameters:

Table F.25 – SITP parameters for master key update, step 8

Parameter	Value	Description
BCF	04 _h	Command “Combined activation/deactivation of security information”
DSI	03 _h	Data structure for combined activation/deactivation
DSH1	FF _h	KeyID not used (no wrapping applied)
DSH2	FF _h	KeyVersion not used (no wrapping applied)
TargetTime	3000000000 _h	set to zero for an instant activation/deactivation
Activated KeyID	00 _h	master key
Activated KeyVersion	n' ^a	according MK'
Deactivated KeyID	00 _h	master key
Deactivated KeyVersion	FF _h	Unused
Option	01 _h	perform a MessageCounter reset
^a If the gateway sets the activated KeyVersion to FFh (unused), the detection of the new master key MK' and the according KeyVersion n' has to be done by the OMS end-device.		

9. The OMS end-device tries to activate KeyID = 00h with version n' and in an atomic operation deactivates KeyID = 00h with version n.
10. In case the gateway does receive a failure, the KeyUpdate is aborted with a permanent failure. The current master key MK is still in use.
11. In case the gateway does receive a successful response, the activation has been successful and gateway and OMS end-device immediately use MK' as the new master key for the next TLS ChannelRequestChannelRequest. If KeyVersion of deactivated MK is >0, MK can be deleted. If KeyVersion is = 0, the initial MK₀ is preserved in the OMS end-device and its associated current MessageCounter is stored persistent in the OMS end-device.

F.4.2.2 Synchronization of state

- 15 In case the gateway does not receive any response after the activation/deactivation or if the update process has been interrupted between Step 7 and Step 10, the gateway may probe, if communication with the OMS end-device is possible by using a message with AFL CMAC using KMAC derived from MK or MK'. If communication using MK is not possible, the probe is done with MK', which must succeed.

F.4.3 Update of OMS end-device certificate and KeyPair

F.4.3.1 General

The OMS end-device holds an individual certificate (MTR_TLS_CRT) with associated KeyPair (MTR_TLS_PUB/MTR_TLS_PRIV) for TLS authentication and (ECDHE) key agreement. The initial OMS end-device certificate and KeyPair is generated and stored in the OMS end-device by the manufacturer. OMS declares the assignment of the MTR_TLS_CRT to the KeyID = 90h and the associated KeyPair to the KeyID = A0h. That means the PublicKey in MTR_TLS_CRT is identical to the PublicKey of the KeyPair (MTR_TLS_PUB). This association is fix and cannot be released, i.e. a kind of rekeying of a certificate is not allowed. For an activation always both components shall be available and valid.

- Each OMS end-device certificate, either stored in the OMS end-device by the manufacturer (initial certificate) or generated by the gateway and transferred to the OMS end-device over a

confidential TLS channel, shall contain a valid date of issue. The date of issue is identified by the “Validity.NotBefore” parameter (see an example in Appendix F.F). The certificate may be self-signed or Issued by a CA (this case is not described here).

The OMS end-device shall be able to manage exact two pairs of credentials - each contains the MTR_TLS_CRT and the associated KeyPair. They are defined as:

- ActiveCredential - the one currently in use
- NewCredential – the one which is intended to replace the active one

There can only be one ActiveCredential and one NewCredential in the OMS end-device. If a new certificate or the KeyPair is transmitted or generated it shall be stored in the NewCredential. During the activation process the NewCredential shall be validated in terms of:

- Syntactical correctness of the data structures
- PublicKey of MTR_TLS_CRT and KeyPair shall be equal
- The date of issue of the NewCredential (contained in MTR_TLS_CRT, see above) shall be newer than the one of the ActiveCredential (if ActiveCredential present).
- The Subject CommonName of the NewCredential shall be the same as the Subject CommonName of the ActiveCredential.
- On initial installation (no ActiveCredential present), the Subject CommonName of the NewCredential shall unambiguously identify the OMS end-device.
- If the NewCredential issuer name is identical with subject name, the certificate is a (self-signed) root Certificate and the signature shall be verified with the PublicKey of the NewCredential. The BasicConstraints extension shall contain $cA = 1$ and $pathLenConstraint = 0$
- If the NewCredential issuer name is different from the subject name, the BasicConstraints extension shall contain $cA = 0$ and the issuer should be a trusted certificate with BasicConstraint $cA = 1$ installed in the OMS end-device.
- If the extension KeyUsage is present in the certificate, the bits for KeyEncipherment and KeyAgreement should be set for MTR_TLS_CRT.

The installation or renewal of the NewCredential can be realized in three different ways:

- Both are generated by the gateway and are transferred to the OMS end-device over a confidential and authentic TLS channel using the SITP-Protocol (see F.4.3.3, with security mode 13)
- If no TLS channel can be established (initial installation) the SITP commands are sent with authentic symmetrical encryption according to OMS security profile B (see F.4.3.3).
- If the OMS end-device generates a public/private KeyPair based on a cryptographically strong random number and the private key is kept confidential in the OMS end-device, the OMS end-device may generate the NewCredential. The gateway retrieves the certificate from the OMS end-device over an authentic channel. The private key does not leave the OMS end-device (see F.4.3.2).

KeyID

According to OMS definition the MTR_TLS_CRT has KeyID = 90h. Associated to MTR_TLS_CRT is an Elliptic-Curve KeyPair (MTR_TLS_PUB, MTR_TLS_PRV) using KeyID = A0h

KeyVersion

Neither the (MTR_TLS_PUB, MTR_TLS_PRV) KeyPair nor the MTR_TLS_CRT are using a KeyVersion. As explained above only the date of issue of the certificate is relevant. Therefore the KeyVersion parameter of SITP shall be set to FF_h and ignored by the OMS end-device.

F.4.3.2 Renewal of KeyPair and certificate by the OMS end-device

The renewal of KeyPair and certificate by the OMS end-device is not supported in this version.

F.4.3.3 Installation of a new KeyPair and new OMS end-device certificate by the gateway

- Transfer of the KeyPair from the gateway to the OMS end-device is done using SITP with the following parameters:

Table F.26 – SITP parameters – Command key transfer

Parameter	Value	Description
BCF	00 _h	Command “Transfer security information”
DSI	11 _h	Data structure for KeyPair transport which contains MTR_TLS_PUB and MTR_TLS_PRV
DSH1	A0 _h	KeyID for KeyPair assigned to MTR_TLS_CRT
DSH2	FF _h	KeyVersion not used
Data structure content		KeyPair ASN.1 DER coded according to RFC5958 chapter 2

- The OMS end-device confirms the transfer of the new key (see Table F.27). In case the gateway does not receive a successful SITP status response, the KeyPair update is aborted with a permanent failure. The ActiveCredential is still in use.

Table F.27 – SITP parameters – Key transfer status ok

Parameter	Value	Description
BCF	80 _h	Response “Transfer security information”
DSI	22 _h	Status response structure
DSH1	A0 _h	KeyID for KeyPair assigned to MTR_TLS_CRT
DSH2	FF _h	KeyVersion not used
Status response byte	00 _h	Successful SITP command

- The OMS end-device stores the new KeyPair into NewCredential, but does not activate it immediately. Activation is done with a separate command after the associated certificate MTR_TLS_CRT is available in the OMS end-device.
- Transfer of the certificate (MTR_TLS_CRT) into the OMS end-device is done using SITP with the following parameters:

Table F.28 – SITP parameters – Command certificate transfer

Parameter	Value	Description
BCF	00 _h	Command “Transfer security information”
DSI	10 _h	Data structure for transporting certificates
DSH1	90 _h	KeyID for MTR_TLS_CRT
DSH2	FF _h	KeyVersion not used
Data structure content		Certificates ASN.1 DER coded according to RFC5280 chapter 4

5. The OMS end-device confirms the transfer of the new key (see Table F.29). In case the gateway does not receive a successful SITP status response, the certificate update is aborted with a permanent failure. The ActiveCredential is still in use.

Table F.29 – SITP parameters – Certificate status ok

Parameter	Value	Description
BCF	80 _h	Response “Transfer security information”
DSI	22 _h	Status response structure
DSH1	90 _h	KeyID for MTR_TLS_CRT
DSH2	FF _h	KeyVersion not used
Status response byte	00 _h	Successful SITP command

- 5 6. Activation of the NewCredential is done using SITP with the following parameters:

Table F.30 – SITP parameters – Command activation

Parameter	Value	Description
BCF	01 _h	Command “Activate security information”
DSI	02 _h	Data structure for activation of security information
DSH1	FF _h	KeyID not used (no wrapping applied)
DSH2	FF _h	KeyVersion not used (no wrapping applied)
TargetTime	3000000000 _h	set to zero for an instant activation/deactivation
KeyID	A0 _h	KeyID for KeyPair assigned to MTR_TLS_CRT
KeyVersion	FF _h	KeyVersion not used - only the NewCredential can be activated

7. If the OMS end-device receives an activation command it shall validate the NewCredential (described in F.4.3.1).

- a. If there are any problems the activation is aborted with a permanent failure. In that case the ActiveCredential is still in use.

If this happen e.g. if the transferred certificate does not match to the transferred KeyPair a SITP status response is sent by the OMS end-device.

Table F.31 – SITP parameters – Activation status with failure

Parameter	Value	Description
BCF	84 _h	Response “Combined activation/deactivation of security information”
DSI	22 _h	Status Response structure
DSH1	FF _h	KeyID not used (no wrapping applied)
DSH2	FF _h	KeyVersion not used (no wrapping applied)
Status response byte	81 _h	Mismatch KeyPair and certificate

- b. If the NewCredential is validated correctly the OMS end-device transfers the NewCredential to the ActiveCredential, deletes the former ActiveCredential and sends a successful SITP response.

Table F.32 – SITP parameters – Activation status with success

Parameter	Value	Description
BCF	84 _h	Response “Combined activation/deactivation of security information”
DSI	22 _h	Status response structure
DSH1	FF _h	KeyID not used (no wrapping applied)
DSH2	FF _h	KeyVersion not used (no wrapping applied)
Status response	00 _h	Successful SITP command

byte		
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8. In case the gateway receives a negative response, the update process is aborted with a permanent failure. The current ActiveCredential is still in use.
9. If the activation has been successful, gateway and OMS end-device use the new ActiveCredential for the next TLS handshake.

F.4.3.4 Synchronization of state

Synchronization could be used in case of doubt about communication state between gateway and OMS end-device e.g. after the reboot of the gateway.

The gateway initiates a TLS handshake (Send TLS ChannelRequest, Receive ClientHello, Send ServerHello + ServerCertificate + ServerKeyExchange + ClientCertificateRequest + ServerHelloDone) and waits for the ClientCertificate (MTR_TLS_CRT). If the OMS end-device sends the MTR_TLS_CRT that corresponds to the activated LMN_TLS_PUB/PRV KeyPair the synchronization has been successful.

F.4.4 Update of gateway certificate (used for TLS ServerTrust)

F.4.4.1 Transfer of gateway certificate

The gateway holds an individual certificate with associated KeyPair for TLS authentication and (ECDHE) key derivation. An OMS end-device should only send data via TLS, after the TLS channel is mutual authenticated, i.e. the presented ServerCertificate matches one of the trust anchors stored in the OMS end-device. As there is no initial gateway certificate in the OMS end-device, the gateway certificate for mutual authentication has to be installed with the SITP protocol using OMS security profile B.

According to OMS definition the self-signed certificate of the gateway (GWLMN_TLS_CRT) is assigned the KeyID 81h or 82h. The public key for use with TLS (GWLMN_TLS_PUB) is derived by the OMS end-device from the gateway certificate GWLMN_TLS_CRT.tbs.subjectPublicKeyInfo. No KeyID is assigned for GWLMN_TLS_PUB.

During the TLS handshake the OMS end-device (TLS client) processes the ServerCertificateMessage sent by the gateway. The OMS end-device shall terminate the TLS handshake, if the presented ServerCertificate (the first in the ServerCertificateMessage) is not equal to one of the certificates stored under KeyID 81h or 82h.

Update of the OMS end-device certificate and gateway certificate is independent from each other and independent from master key update MK.

After changing the OMS end-device certificate or gateway certificate, the TLS session keys are still valid. The new certificates come into effect, when the next full TLS handshake is started.

For installation of the initial GWLMN_TLS_CRT no TLS channel can be used. In this case the following SITP commands are sent with symmetrical encryption using the current master key MK according to OMS security profile B. For update of subsequent GWLMN_TLS_CRT a TLS channel shall be used. In this case the following SITP commands are sent according to OMS security profile C (with CF. security mode 13).

NOTE: Typically, the MK0 is active when transferring the initial GWLMN_TLS_CRT.

1. Transfer of the certificate (GWLMN_TLS_CRT) into the OMS end-device is done using SITP with the following parameters:

Table F.33 – SITP parameters – Transfer certificate

Parameter	Value	Description
BCF	00 _h	Command "Transfer security information"

DSI	10 _h	Data structure for transporting certificates
DSH1	81 _h or 82 _h	KeyID for GWLMN_TLS_CRT
DSH2	FF _h	KeyVersion not used for certificates
Data structure content		Certificates ASN.1 DER coded according to RFC5280 chapter 4

2. In case the gateway does not receive a successful SITP response, the certificate update is aborted with a permanent failure. The currently active GWLMN_TLS_CRT is still in use.
3. Activation of the GWLMN_TLS_CRT is done using SITP with the following parameters:

Table F.34 – SITP parameters – Command activation

Parameter	Value	Description
BCF	01 _h	Command “Activate security information”
DSI	02 _h	Data structure for activation of security information
DSH1	FF _h	KeyID not used (no wrapping applied)
DSH2	FF _h	KeyVersion not used (no wrapping applied)
TargetTime	3000000000 _h	set to zero for an instant activation/deactivation
KeyID	81 _h or 82 _h	KeyID for KeyPair assigned to MTR_TLS_CRT
KeyVersion	FF _h	KeyVersion not used - only the NewCredential can be activated

4. For an activation of the new GWLMN_TLS_CRT the validations (syntactical correctness and date-of-issue check) shall be done successfully. If there are any problems the activation is aborted and a SITP status response (containing the error) is sent by the OMS end-device.
5. In case the activation was successful the new certificate replaces the current one with the identical KeyID in the OMS end-device (e.g. both have 81_h). Gateway and OMS end-device use the certificate for verification of the ServerCertificate in the next (full) TLS handshake. A positive SITP status response is sent by the OMS end-device.

F.4.4.2 Synchronization of state

Upon TLS ChannelRequest from gateway to OMS end-device, the OMS end-device sends ClientHello to the gateway and receives a ServerCertificateMessage containing the GWLMN_TLS_CRT of the gateway. If the first certificate of the ServerCertificateMessage is not found in the OMS end-device under KeyID 81h-82h, no communication with the gateway is allowed and a reset to MK₀ is required.

Alternative: The gateway uses SITP with security profile B to read the certificates of KeyID = 81h or 82h with the following parameters:

Table F.35 – SITP parameters – Command “Get security information”

Parameter	Value	Description
BCF	06h	Command “Get security information”
DSI	00h	Empty structure used for requesting security information
DSH1	81h or 82h	KeyID for GWLMN_TLS_CRT
DSH2	FFh	KeyVersion not used for certificates

The response will deliver the currently active certificates in the OMS end-device, e.g.:

Table F.36 – SITP parameters – Response ‘Get security information’

Parameter	Value	Description
BCF	86 _h	Response “Get security information”
DSI	10 _h	Data structure for transporting certificates
DSH1	81 _h	KeyID for GWLMN_TLS_CRT
DSH2	FF _h	KeyVersion not used for certificates
Data structure content		Certificates ASN.1 DER coded according to RFC5280 chapter 4

F.4.5 Reset to MK₀

Using physical access to the OMS end-device, the master key may be restored to MK₀ (KeyID = 0, KeyVersion = 0). Such a reset of MK₀ shall lead to the deletion of all other master keys with KeyID = 0 and KeyVersion > 0 and any GWLMN_TLS_CRT with KeyID 81h-82h. Restoring the master key MK₀ also restores the associated initial MessageCounter C_{M0} (see [OMS-S2], 9.3.2.1.2).

NOTE: As the initial MessageCounter C_{M0} is not resettable, the MK₀ should be replaced by a new MK after establishing a TLS-communication to the OMS end-device (see also [OMS-S2], 9.3.2.1.2).

F.5 TLS-Certificates for OMS end-devices

An OMS end-device shall be able to handle (i.e. process and store) X.509v3 Certificates (according to RFC5280) DER-encoded. Details are shown in Table F.37. See also an example in

5 Appendix F.F.

NOTE: The [BSITR03109-1] can request additional parameters than shown here.

Table F.37 – TLS certificate details

Parameter	Description
Certificate Size	Certificate Size up to 500 Bytes minimum sufficient for EC key sizes up to brainpoolP512r1 without further certificate extensions
PublicKey	Uncompressed PublicKey of Type id-ecPublicKey according to RFC5480 Chap. 2.2
Curves ^a	<ul style="list-style-type: none"> Required: brainpoolP256r1, NIST P256 Recommended: brainpoolP384r1, brainpoolP512r1, NIST P384
Digest Algorithm ^a	<ul style="list-style-type: none"> Required: SHA256 Recommended: SHA-SHA384, SHA-512
Signature Algorithm ^a	Required: ECDSA
SerialNumber	Length 4..20 bytes
BasicConstraintsExtension (critical)	<ul style="list-style-type: none"> cA = 1 pathLenConstraint = 0
KeyUsage Extension (critical)	DigitalSignature Required
Distinguished Name (DN)	<ul style="list-style-type: none"> Selfsigned, i.e. issuerDN = subjectDN commonName (case Insensitive): accept 1..64 characters of ASN.1 type "Printable string". Currently national format empty or according to DIN43863-5:2012-04 without spaces, appended by .mtr/.MTR. Example "7mfc0100001234.mtr". serialNumber attribute (OID 1.3.6.1.4.1.1466.115.121.1.44) accept Sequence-Number of certificate Currently no other DN-attributes are present, but certificate should not be rejected, if they are present. For a transition period also certificates with empty issuerDN, subjectDN (no attributes) must be accepted.
Validity	<ul style="list-style-type: none"> notBefore: UTCTIME and GENERALIZED TIME notAfter: UTCTIME and GENERALIZED TIME
Other extensions	Currently no other extensions are present, but certificate should not be rejected, if they are present.
^a Deprecation and announcement of supported Algorithms and key sizes follows annually updated BSI TR-03116-3.	

Appendix F.A (normative): Security Information Transfer Protocol (SITP)

F.A.1 Introduction

This specific application protocol is intended for all kind of security information handling and management of security relevant services in a metering system. The main use case is to update key information. Handling of security information requires bidirectional access to the OMS end-device. This annex extends the SITP in [EN13757-7:2018], Annex A by services for asymmetric crypto methods such as TLS.

In detail the SITP provides the following services:

- Transfer security information (to OMS end-device)
- Activate security information
- Deactivate security information
- Combined activation/deactivation of security information
- Generate security information
- Get security information (from OMS end-device)
- Get list of all security information (from OMS end-device)
- Get list of active security information (from OMS end-device)
- Transfer end to end secured application data

The usage of this upper layer protocol is introduced with a specific sub-range of CI-fields. It is therefore independent from the M-Bus application layer protocol.

Two examples which show the usage of the SITP are provided in the CEN TR 17167; F.7.

F.A.2 SITP services

F.A.2.1 Transfer security information

The purpose of this service is to transfer security information (e.g. security keys) from a communication partner to an OMS end-device in a secure way. The security information itself is often generated by a backend system and not by the communication partner. This service provides, in addition to security service from the lower layers a wrapping mechanism. The wrapping mechanism ensures a strong binding between the key and the set of data specifying the use of the key. The transferred security information contains a target time that gives an absolute or relative timestamp of usage. The target time may as well be given by the separate activation or deactivation service described below.

NOTE: This service gives no advice on how to handle existing security information when new information is received. Such handling is a part of the security policy of the whole system.

F.A.2.2 Activate security information

This service shall not be used for the OMS. Use “Combined activation/deactivation of security information” instead.

F.A.2.3 Deactivate security information

This service shall not be used for the OMS. Use “Combined activation/deactivation of security information” instead.

F.A.2.4 Destroy security information

The purpose of this service is to destroy (and delete) specific security information formerly used by the OMS end-device. It avoids possible security issues with old security information material. The OMS end-device shall not allow destroying active security information.

F.A.2.5 Combined activation/deactivation of security information

The purpose of this service is to combine the action of activation and deactivation. It is needed in case of a “message counter sharing” between an old and a new key. As a new key requires a reset of the MessageCounter the old key cannot be used with this counter value zero anymore. To avoid security issues in that situation both actions have to take place simultaneously.

F.A.2.6 Generate security information

The purpose of this service is to instruct the OMS end-device to create security information (e.g. a new key pair, a certificate or a certificate signing request “CSR”). The security information can be requested by a separate command.

F.A.2.7 Get security information

The purpose of this service is to get single security information from the OMS end-device. This can be used to ask for a certificate, a CSR, the public key or applied key id, key version of this OMS end-device.

F.A.2.8 Get list of all key information

The purpose of this service is to get a list of all used key id and the available key versions stored in an OMS end-device.

F.A.2.9 Get list of active key information

The purpose of this service is to get a list of all active key id, their key versions and the key counter, which are currently valid.

F.A.2.10 Transfer end to end secured application data

The purpose of this service is to securely transfer an application layer protocol message between a remote command initiator and an OMS end-device. The application data are integrity and maybe privacy protected by usage of a key which is only known to the OMS end-device and the command initiator. An intermediary will not be able to alter this message as the end parties are able to detect this modification. This allows an end to end secured transfer of critical commands or critical data.

F.A.3 CI-fields

Commands and responses for the Security Information Transfer Protocol (SITP) are identified by specific values of the CI-field. The applicable CI-fields are listed in the Table A.1:

Table F.A.1 — SITP CI-fields

CI-field	Designation	Header	Remarks
C3h	Command to device	Long	Security Information Transfer Protocol
C4h	Response from device	Short	Security Information Transfer Protocol
C5h	Response from device	Long	Security Information Transfer Protocol

F.A.4 SITP structure

The SITP structure is implemented as APL data (see EN 13757-7:2018, Table 10 and [OMS-S2], 8.1). It provides the possibility to transfer multiple blocks of security information in one message. Each block contains a single SITP command or SITP response (see Table F.A.3). The block is identified by a block ID which shall be used for a correlation between commands and responses. The function of each block is specified by the block control field. Each command and each response shall have a principal structure as depicted in the table below:

Table F.A.2 - Internal block structure of SITP

Block length (BL)	Block ID field (BID)	Block control field (BCF)	Block parameters
2 bytes	1 byte	1 byte	<i>n</i> bytes

The elements are:

Block length: BL, Length of this block including BID, BCF and all block parameters (LSB first). BL = 0 signals no more blocks to follow

Block ID field: BID, Identifier of this block, start with first block BID = 0

Block control field: BCF, coding of the usage of this command or response, see Table F.A.5

Block parameters: Variable parameters depending on the BCF, see F.A.6

Block length is a multi-byte field. Block parameters may also contain multi byte fields. If nothing else is declared then multi byte fields shall be transmitted with the least significant byte first (little endian).

F.A.5 Block control field

The block control field gives the function of the block and shall have one of the listed values for command or response as shown in the table below. There is a fixed relation between a certain command and its assigned response, e.g. command 03_h requires a response with 83_h.

Table F.A.3 - Block control field

Command	Applicable DSI	Function	Response	Applicable DSI
00 _h	01 _h , 10 _h , 11 _h	Transfer security information (see F.A.2.1)	80 _h	22 _h
01 _h	02 _h	Activate security information (see F.A.2.2)	81 _h	22 _h
02 _h	02 _h	Deactivate security information (see F.A.2.3)	82 _h	22 _h
03 _h	02 _h	Destroy security information (see F.A.2.4)	83 _h	22 _h
04 _h	03 _h	Combined activation/deactivation of security information (see F.A.2.5)	84 _h	22 _h
05 _h	02 _h	Generate security information (see F.A.2.6)	85 _h	22 _h
06 _h	00 _h	Get security information (see F.A.2.7)	86 _h	10 _h , 12 _h , 13 _h , 22 _h

Command	Applicable DSI	Function	Response	Applicable DSI
07 _h	00 _h	Get list of all key information (see F.A.2.8)	87 _h	20 _h , 22 _h ,
08 _h	00 _h	Get list of active key information (see F.A.2.9)	88 _h	21 _h , 22 _h ,
09 _h ^b	00 _h	Get list of active keys and key counter information (see F.A.2.9)	89 _h	22 _h , 23 _h ^b ,
0A _h to 1F _h		Reserved for future use	8A _h to 9F _h	
20 _h	30 _h to 37 _h	Transfer end to end secured application data (see F.A.2.10)	A0 _h	22 _h , 30 _h to 37 _h ,
21 _h to 6F _h		Reserved for future use	A1 _h to EF _h	
70 _h to 7F _h	F0 _h to FF _h ^a	Manufacturer specific usage	F0 _h to FF _h	F0 _h to FF _h ^a
^a All other DSI can be used as well				
^b These values are not yet covered by [EN13757-7:2018]				

F.A.6 Block parameters

The structure of the block parameter is defined in Table F.A.4. It is applied for all SITP commands and responses. In case of status responses this structure is also used but the data structure content will be different (see Table F.A.5).

- 5 The structure is designed to transport applicative data in the “Data structure content”. The other fields serve as organisational and introductive information. The RecipientID and the DSH shall have the same values in a command and the respective response. This is essential for a correct assignment of the response by the command initiator.

Table F.A.4 - Block parameter structure

RecipientID	Data structure identifier (DSI)	Data structure header (DSH)		Data structure content
		DSH1	DSH2	
1 byte	1 byte	1 byte	1 byte	<i>n</i> bytes

- 10 Recipient ID: Identifier of the application. Set to 00_h if no dedicated application is addressed.
- NOTE:** For example to address the gateway function in a combined electricity meter. In the TCP world this would be the port number.
- 15 Data structure identifier: DSI, identifying the applied data structure declared in Table F.A.5
- Data structure header: DSH introduces the data structure. Its meaning depends on the data structure identifier. DSH1 and DSH2 are autonomous 1-byte fields as shown in Table F.A.5.
- Data structure content: Structured security information according to F.A.7

F.A.7 Overview about data structures / mechanisms

This chapter informs about the data structures, the applied mechanisms and the usage of the DSH. The “data structure content” as shown in Table F.A.4 is provided in the dedicated sub clauses (e.g. data structure 01_h) or defined by the data structure itself (e.g. CSR according to RFC4211 chapter 3). In this case no further explanations about the data structure content are given in concerning section.

There are two different groups of data structures defined which can be seen in Table F.A.5. One group is mainly intended for handling and management of security information and their relevant services (DSI 00_h to 2F_h, see F.A.8). The other group is defined for the use case of secured application data transfer (DSI 3x_h, see F.A.9). Table F.A.5 shows the list of data structures / mechanisms, the usage of DSH and the applicable BCF values. Details about the DSH values can be found below the table.

Table F.A.5 – List of SITP data structures / mechanisms

Data Structure Identifier (DSI)	Data structure / Mechanism	Data structure header (DSH)	
		DSH1	DSH2
00 _h	Empty structure / no wrapping (see [EN13757-7:2018], Annex A, A.8.2)	KeyID	KeyVersion
01 _h	Wrapping AES128 according NIST SP 800-38F type KWP (see [EN13757-7:2018], Annex A, A.8.3)	WrapperKey ID	WrapperKey Version
02 _h	Wrapping AES128 according NIST SP 800-38F type KWP (see [EN13757-7:2018], Annex A, A.8.4)	WrapperKey ID	WrapperKey Version
03 _h	Wrapping AES128 according NIST SP 800-38F type KWP (see [EN13757-7:2018], Annex A, A.8.5)	WrapperKey ID	WrapperKey Version
04 _h to 0F _h	Reserved for future use	n/a	n/a
10 _h	Certificates ASN.1 DER coded according to RFC5280 chapter 4	KeyID	unused ^a
11 _h	KeyPair ASN.1 DER coded according to RFC5958 chapter 2	KeyID	unused ^a
12 _h	Public keys ASN.1 DER coded according to subject public key info (RFC5480 chapter 2) ^c	KeyID	unused ^a
13 _h	CSR according to RFC4211 chapter 3	KeyID	unused ^a
14 _h to 1F _h	Reserved for asymmetric technologies	n/a	n/a
20 _h	All keys structure wrapped with AES128 according NIST SP 800-38F type KWP (see [EN13757-7:2018], Annex A, A.8.6.)	WrapperKey ID	WrapperKey Version
21 _h	Active keys structure ([EN13757-7:2018], Annex A, A.8.7)	unused ^a	unused ^a
22 _h	Status response structure (see [EN13757-7:2018], Annex A, A.8.8)	KeyID ^b	KeyVersion ^b
23 _h ^d	Active key counter structure (F.A.8)	unused ^a	unused ^a
24 _h to 2F _h	Reserved for future use	n/a	n/a
30 _h	Wrapping AES128 according to NIST SP 800-38F type KWP ([EN13757-7:2018], Annex A, A.9.2)	WrapperKey ID	WrapperKey Version
31 _h	Authentication with HMAC SHA256 (see [EN13757-7:2018], Annex A, A.9.3)	WrapperKey ID	WrapperKey Version
32 _h	Authentication with AES128 CMAC (8 byte MAC, see [EN13757-7:2018], Annex A, A.9.4)	WrapperKey ID	WrapperKey Version
33 _h	Authentication with AES128 CMAC (16 byte MAC, see [EN13757-7:2018], Annex A, A.9.4)	WrapperKey ID	WrapperKey Version
34 _h	Authenticated encryption with AES128-GCM (see [EN13757-7:2018], Annex A, A.9.5)	WrapperKey ID	WrapperKey Version
35 _h	Authentication with AES128 GMAC (see [EN13757-7:2018], Annex A, A.9.6)	WrapperKey ID	WrapperKey Version
36 _h	Authenticated encryption with AES128 CCM (8 byte MAC, see [EN13757-7:2018], Annex A, A.9.7)	WrapperKey ID	WrapperKey Version
37 _h	Authenticated encryption with AES128-CCM (16 byte MAC, see [EN13757-7:2018], Annex A, A.9.7)	WrapperKey ID	WrapperKey Version
38 _h to EF _h	Reserved for future use	N/A	NA
F0 _h to FF _h	Manufacturer specific	N/A	N/A
^a Set unused fields to FF _h ^b The same DSH value than used in command ^c This chapter extends RFC5280 4.1.2.7 ^d This value are not yet covered by [EN13757-7:2018]			

Contents of data structure header (DSH):

WrapperKeyID: ID of the key used for wrapping or authentication according the KeyID definition. If not used (no wrapping applied) set to FFh. Note that only KeyID in the range from 00h to 0Fh are used for TPL security (see Table 24). WrapperKeys used for the APL security shall use a KeyID larger than 0Fh.

WrapperKeyVersion: Version of the key used for wrapping. If not used set to FFh

KeyID: The ID of the intended key. If not used set to FFh.
OMS declares specific meaning for certain KeyID.

Table F.A.6 - Predefined OMS KeyID

KeyID	Type of Security information
00 _h to 7F _h	see [OMS-S2]; 9.2.2
80 _h	n/a
81 _h to 82 _h ^a	Trust anchor certificate (GWLMN_TLS_CRT)
83 _h to 8F _h	Reserved for Trust anchor certificate
90 _h	Device specific TLS server authentication certificate (MTR_TLS_CRT)
91 _h	Application data signature certificate (MTR_SIG_CRT)
92 _h to 9F _h	n/a
A0 _h	KeyPair assigned to MTR_TLS_CRT
A1 _h	KeyPair assigned to MTR_SIG_CRT
A2 _h to AF _h	n/a
B0 _h to FF _h	see [OMS-S2], 9.2.2
^a A second trust anchor may be necessary for a fail-safe transfer of authority from one SMGW to another SMGW with a direct trust concept.	

KeyVersion: The version of the intended KeyID. If not used set to FFh.

F.A.8 Data structures for security information

The OMS end-device shall apply a data structures for security information according [EN13757-7:2018], Annex A, A.8.

For the Data structure 22_h (status response) the values of [EN13757-7:2018], Annex A, table A.13 are extended by following table.

Table F.A.7 – Extension of status response byte definition

Status response byte	Explanation
09 _h	This SITP block was not executed due to an error in another SITP block
0A _h to 0F _h	Reserved for status information (no error)
...	...
11 _h	Unknown or invalid command (BCF error)
12 _h to 13 _h	Reserved for data structure header errors
...	...
80 _h	Unspecified certificate error
81 _h	Mismatch KeyPair and certificate
82 _h to 8F _h	Reserved for asymmetric usage

Data structure 23_h

The Data structure 23_h extend the SITP in [EN13757-7:2018], with a missing functionality. This data structure is used to report the current KeyCounter values used in 0. The KeyID, the KeyVersions and the stored KeyCounter will be transferred, not the keys itself. This structure provides the possibility to ask for the current KeyCounter value per KeyID. It is not wrapped/protected as the KeyID and KeyCounter are no secret information. The length of the whole structure is variable, as the number of used KeyID is not defined as shown in Table F.A.8 (the maximum length is 256 * 6 bytes).

Table F.A.8 – Active key counter structure 23_h

KeyID	KeyVersion	KeyCounter
1 byte	1 byte	4 byte
...
1 byte	1 byte	4 byte

KeyID: 1 byte value of the reported key id. Only key id which has an activekey verison shall be given here.

KeyVersion: 1 byte value for the active key version of this key id. If no key versioning used set to FFh

KeyCounter: 4 byte KeyCounter (last stored one) of the respective key id with LSB first.

F.A.9 Data structures for secured application data

For the secured transfer of application protocols with critical commands or critical data the OMS end-device may apply data structures according to [EN13757-7:2018], Annex A, A.9:

- Data Structure 30h – AES key wrap (encryption and authentication) or
- Data Structure 32h and 33h – CMAC (authentication only).

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The use of other data structures for secured application data is optional and will not be checked by the OMS-CT.

Appendix F.B (informative): Examples for the usage of AFL and TLS

This chapter shows several examples of messages using AFL and TLS.

Table F.B.1 – Simple secured unfragmented M-Bus message

DLL header (EN 13757-4): Len,MsgType Addr	Long/Short Transport App header CI = 72h, ADDR, ACC, ST, CF = MODE05	Application data [OMS-S2] M-Bus DIF/VIFs
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Table F.B.2 – Non-fragmented authenticated M-Bus message

DLL header (EN 13757-4): Len,MsgType Addr	Extended Link Layer CI = 8Ch CC, ACC [OMS-S2]	Authentication and Fragmentation Layer CI = 90h AFLLen, AFLFlags, CTR, MAC	Long/Short Transport header CI = 5Bh, ADDR, ACC, ST, CF = MODE07	Application data [OMS-S2]
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Table F.B.3 – Fragmented authenticated application message (1st fragment)

DLL header (EN 13757-4): Len,MsgType, Addr	Extended Link Layer CI = 8Eh, CC, ACC, ADDR [OMS-S2]	Authentication and Fragmentation Layer CI = 90h, AFLLen, AFLFlags(MF = 1), TLSSDULen, CTR	Long/Short Transport header CI = 7Ah, ACC, ST, CF = MODE07	Application data [OMS-S2]
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Table F.B.4 – Fragmented authenticated application message (last fragment)

DLL header (EN 13757-4): Len,MsgType, Addr	Extended Link Layer CI = 8Eh, CC, ACC, ADDR [OMS-S2]	Authentication and Fragmentation Layer CI = 90h, AFLLen, AFLFlags(MF = 0), MAC	Application data [OMS-S2]
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Table F.B.5 – Unfragmented TLS secured application command message

DLL header (EN 13757-4): Len,MsgType, Addr	Extended Link Layer CI = 8Ch CC, ACC [OMS-S2]	Long/Short Transport header CI = 5Bh, ADDR, ACC, ST, CF = MODE13 [OMS-S2]	TLS record (includes application data)		
			TLS HDR, IV	M-Bus application data [OMS-S2]	HMAC/ Padding

Appendix F.C (informative): Master key renewal

This Appendix shows the key renewal procedure as described in F.4.2.

Figure F.C.1 – Key renewal procedure – Overview

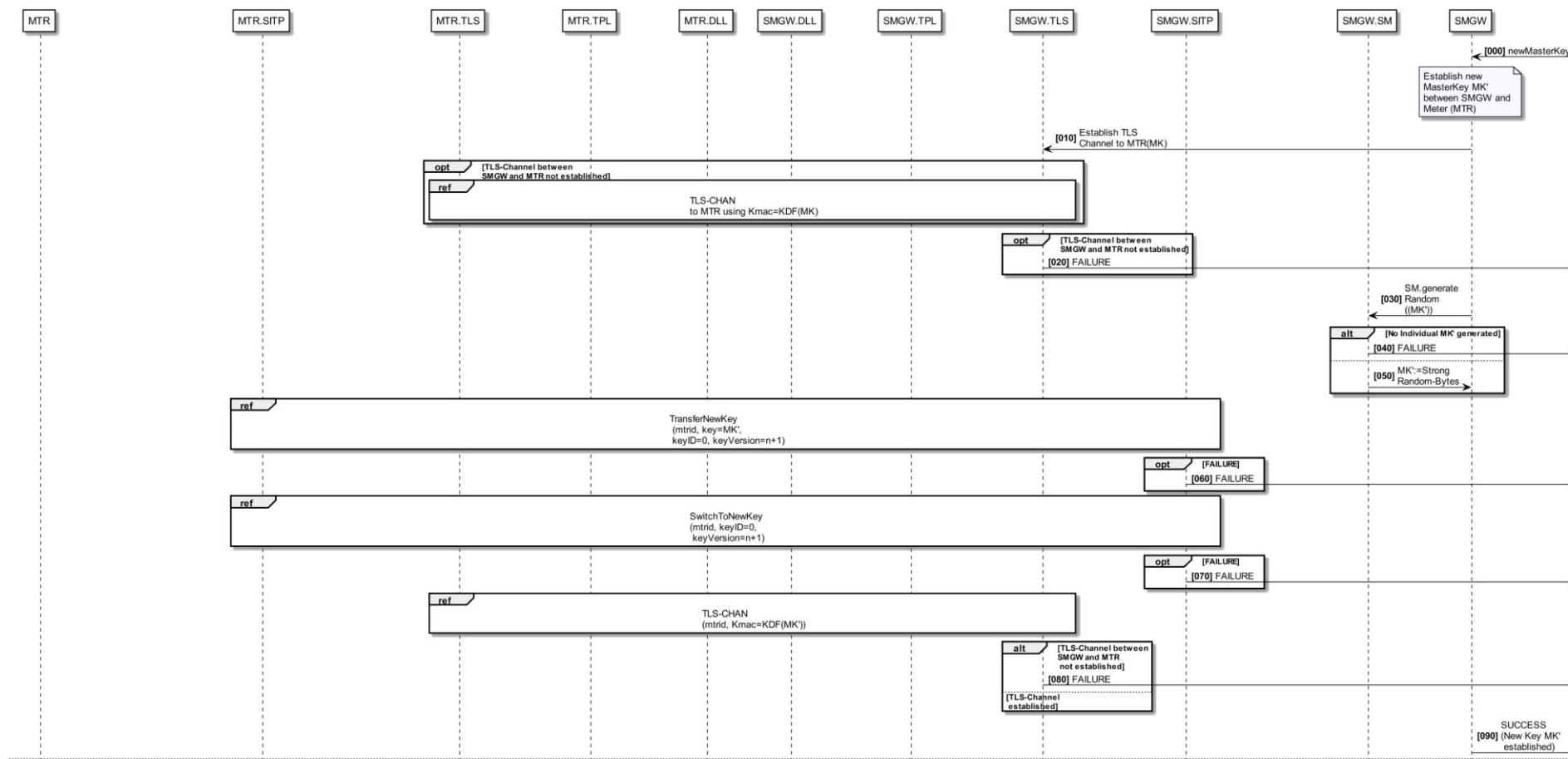


Figure F.C.2 – Key renewal procedure – Key transfer

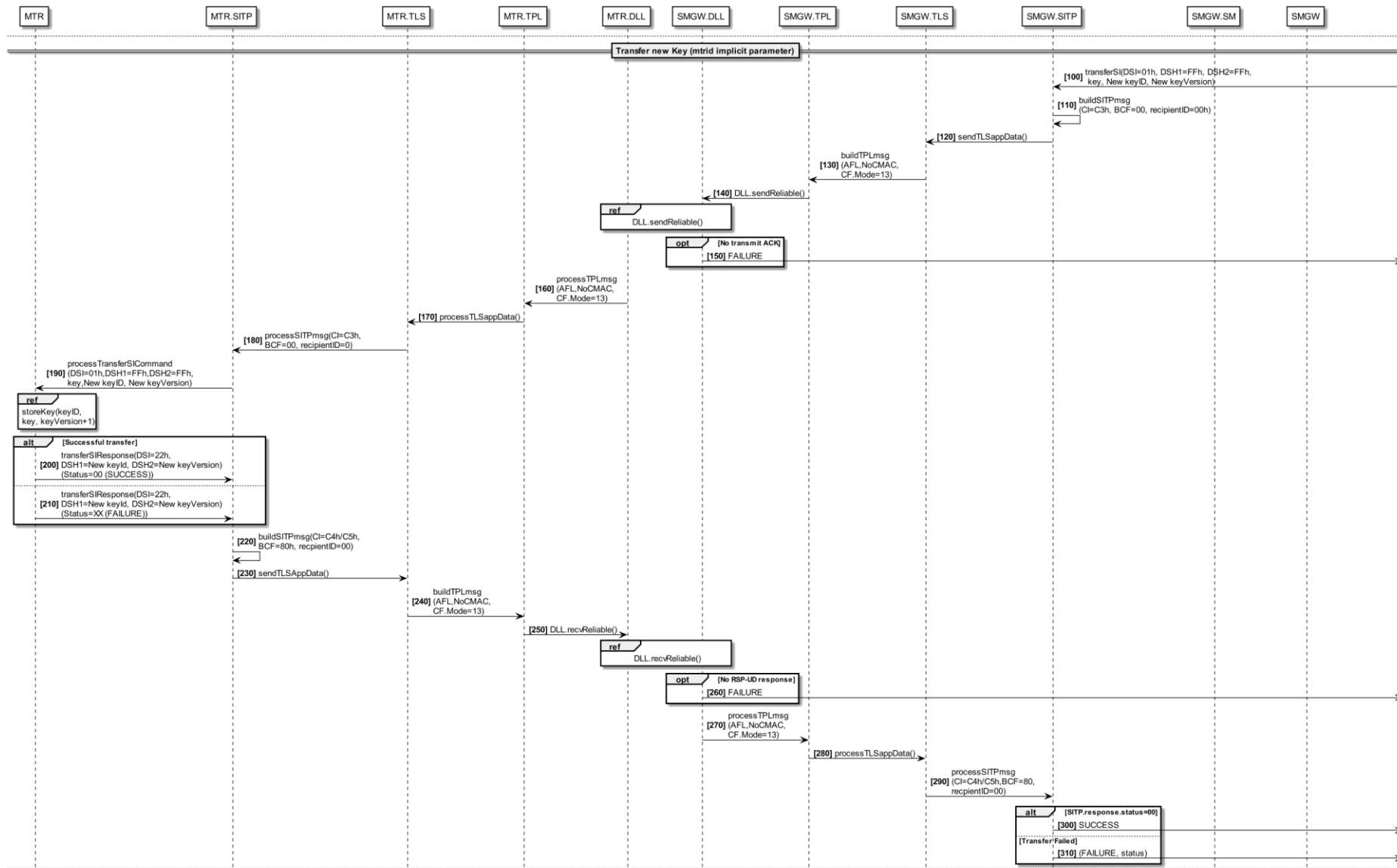


Figure F.C.3 – Key renewal procedure – Key activation/deactivation

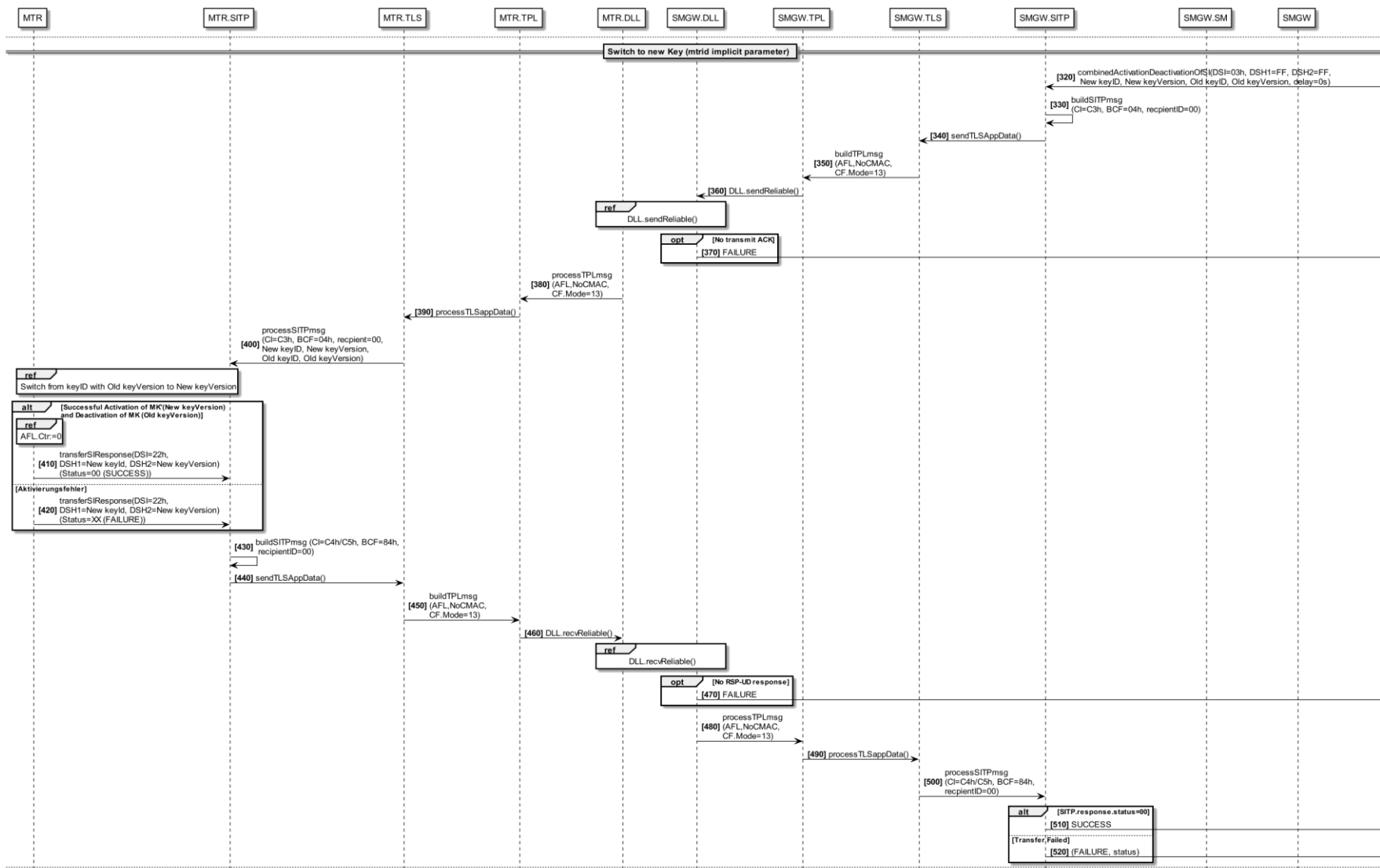
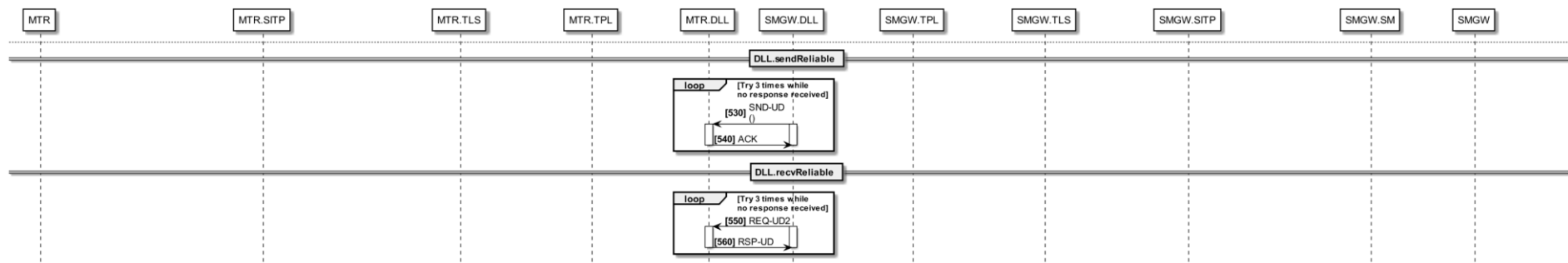


Figure F.C.4 – Key renewal procedure – Send and receive of datagrams



MK: Current MasterKey (Key Version n)
MK': New MasterKey (Key Version n+1)
SITP: Security Information Transport Protocol
AFL: Authentication and Fragmentation Layer
DLL: Data Link Layer
TPL: Transport Layer
TLS: Transport Layer Security
MTR: Meter
SMGW: Smart Meter Gateway
SM: Security Module in SMGW
KDF: Key Derivation Function
mtrid: Meter-Address
gwid: Gateway-Address

Appendix F.D (informative): Message examples of TLS

F.D.1 General

The colours are applied for the distinction of the different layers. They are used according to figure 1 of the superior document [OMS-S2].

- 5 The column "Offset" shows the length of the application data calculated from the beginning of the application layer.

NOTE 1: In this example all messages from the OMS end-device use a short header (no radio adapter on the OMS end-device).

- 10 **NOTE 2:** In example of radio frames are preamble, synchronisation word and CRC-fields not presented to get a better readability.

F.D.2 TLS ChannelRequest (from server to client)

F.D.2.1 Example for wireless M-Bus

Table F.D.1 – TLS-ChannelRequest

Layer	0 (first byte)	1	2	3
DLL	L = nn	C = 53h/73h (SND-UD)	MFCT_0	MFCT_1
DLL	ID_0(GW)	ID_1(GW)	ID_2(GW)	ID_3(GW)
DLL	VER(GW)	DEVTYPE(31h)		
ELL	CI = 8Ch(ELL)	CC	ACC	
AFL	CI = 90(AFL)	AFL.AFLL (0Fh)	AFL.FCL(01h) LSB	AFL.FCL(2Ch) MSB MF = 0,CTR, MCL, MAC
AFL	AFL.MCL(25h), CMAC8,CTR	AFL.MCR[7..0]	AFL.MCR[15..8]	AFL.MCR[23..16]
AFL	AFL.MCR[31..24]	AFL.MAC	AFL.MAC	AFL.MAC
AFL	AFL.MAC	AFL.MAC	AFL.MAC	AFL.MAC
AFL	AFL.MAC			
TPL	CI = 5Fh	ID_0(MTR)	ID_1(MTR)	ID_2(MTR)
TPL	ID_3(MTR)	MFCT_0(MTR)	MFCT_1(MTR)	VER(MTR)
TPL	DEVTYPE(MTR)	ACC	ST	CF[7..0]
TPL	CF[15..8](Mode13)	CFE[7..0] (TLSCHAN)		
TPL	ContentType (00h, ChannelRequest)	Reserved (00h)	Reserved (00h)	TLSSDULen MSB (00h)
TPL	TLSSDULen LSB (00h)			

- 15 The OMS end-device answers with ACK (no data) and waits for a REQ-UD2 from the gateway.

F.D.2.2 Example for wired M-Bus

Table F.D.2 – TLS ChannelRequest

Layer	0 (first byte)	1	2	3
DLL	68h	L = nn	L = nn	68h
DLL	C = 53h/73h (SND-UD)	PAddr (Primary M-Bus addr.)		
AFL	CI = 90h (AFL)	AFL.AFLL (0Fh)	AFL.FCL (00h) LSB FID = 0 (no fragments)	AFL.FCL (2Ch) MSB MF = 0 MCLP = 1 MLP = 0 MCRP = 1 MACP = 1
AFL	AFL.MCL (25h) MLMP = 0 MCMP = 1 AT = CMAC AES 128 Trunc.8 bytes	AFL.MCR [7..0] (xxh)	AFL.MCR [15..8] (xxh)	AFL.MCR [23..16] (xxh)
AFL	AFL.MCR [31..32] (xxh)	AFL.MAC [63..56] (xxh)	AFL.MAC [55..48] (xxh)	AFL.MAC [47..40] (xxh)
AFL	AFL.MAC [39..32] (xxh)	AFL.MAC [31..24] (xxh)	AFL.MAC [23..16] (xxh)	AFL.MAC [15..8] (xxh)
AFL	AFL.MAC [7..0] (xxh)			
TPL	CI = 5Fh	ID_0(MTR)	ID_1(MTR)	ID_2(MTR)
TPL	ID_3(MTR)	MFCT_0(MTR)	MFCT_1(MTR)	VER(MTR)
TPL	DEVTYPE(MTR)	ACC	STS	CF[7..0]
TPL	CF[15..8](MODE 13)	CFE[7..0](TLSCHAN)		
TPL	ContentType (00h, Application Data)	Reserved (00h)	Reserved (00h)	TLSSDULen (MSB, 00h)
TPL	TLSSDULen (LSB, 00h)			
DLL	Chksum	0x16		

F.D.3 First message flight – ClientHello (from client to server)

Table F.D.3 – First message flight (TLS session initiate - Full handshake)

Layer	0 (first byte)	1	2	3
DLL	L = nn	C = 08h/28h (RSP-UD)	MFCT_0(MTR)	MFCT_1(MTR)
DLL	ID_0(MTR)	ID_1(MTR)	ID_2(MTR)	ID_3(MTR)
DLL	VER(MTR)	DEVTYPE(MTR)		
ELL	CI = 8C(ELL)	CC	ACC	
AFL	CI = 90h (AFL)	AFL.AFLL (0Fh)	AFL.FCL [7..0] (00h) FID = 0 (no fragments)	AFL.FCL [15..8] (2Ch) MF = 0 MCLP = 1 MLP = 0 MCRP = 1 MACP = 1
AFL	AFL.MCL (25h) MLMP = 0 MCMP = 1 AT = CMAC AES 128 Trunc.8 bytes	AFL.MCR [7..0] (xxh)	AFL.MCR [15..8] (xxh)	AFL.MCR [23..16] (xxh)
AFL	AFL.MCR [31..32] (xxh)	AFL.MAC [63..56] (xxh)	AFL.MAC [55..48] (xxh)	AFL.MAC [47..40] (xxh)
AFL	AFL.MAC [39..32] (xxh)	AFL.MAC [31..24] (xxh)	AFL.MAC [23..16] (xxh)	AFL.MAC [15..8] (xxh)
AFL	AFL.MAC [7..0] (xxh)			
TPL	CI = 9Eh	ACC	STS	CF[7..0]
TPL	CF[15..8] (MODE 13,)	CFE[7..0] (TLSPROT)		
TPL	ContentType (16h, Handshake)	ProtoMajor (03h) TLS1.2	ProtoMinor(03h) TLS1.2	TLSSDULen (MSB, 00h)
TPL	TLSSDULen (LSB, 4Ch)	HSMmsgType (01h ClientHello)	HSMmsgLength (MSB, 00h)	HSMmsgLength (00h)
TPL	HSMmsgLength (LSB)	HSClientVersionMajor (03h)	HSClientVersionMinor (03h) TLS1.2	HSClientRandom 32
TPL	HSClientRandom 31	HSClientRandom	HSClientRandom	HSClientRandom
TPL	HSClientRandom 27	HSClientRandom	HSClientRandom	HSClientRandom
TPL	HSClientRandom 23	HSClientRandom	HSClientRandom	HSClientRandom
TPL	HSClientRandom 19	HSClientRandom	HSClientRandom	HSClientRandom
TPL	HSClientRandom 15	HSClientRandom	HSClientRandom	HSClientRandom
TPL	HSClientRandom 11	HSClientRandom	HSClientRandom	HSClientRandom
TPL	HSClientRandom 7	HSClientRandom	HSClientRandom	HSClientRandom
TPL	HSClientRandom 3	HSClientRandom 2	HSClientRandom 1	HSSessionIDLength (00h No Resume) ¹
TPL	CipherSuitesLength (MSB, 00h)	CipherSuitesLength (LSB, 08h)	CipherSuite Prio1 (MSB, C0h) TLS_ECDHE_ECDSA_WITH_AES_128_CBC_SHA256	CipherSuitePrio1 (LSB, 23h) TLS_ECDHE_ECDSA_WITH_AES_128_CBC_SHA256

¹ **NOTE:** If no TLS session resumption is requested, the SessionID length is 00h. With TLS session resumption the maximum length is 32 bytes.

Layer	0 (first byte)	1	2	3
TPL	CipherSuitePrio2 (MSB, C0h) TLS_ECDHE_ECDSA_WITH_AES_256_CBC_SHA384	CipherSuitePrio2 (LSB, 24h) TLS_ECDHE_ECDSA_WITH_AES_256_CBC_SHA384	CipherSuitePrio3 (MSB, C0h) TLS_ECDHE_ECDSA_WITH_AES_128_GCM_SHA256	CipherSuitePrio3 (LSB, 2Bh) TLS_ECDHE_ECDSA_WITH_AES_128_GCM_SHA256
TPL	CipherSuitePrio4 (MSB, C0h) TLS_ECDHE_ECDSA_WITH_AES_256_GCM_SHA384	CipherSuitePrio4 (LSB, 2Ch) TLS_ECDHE_ECDSA_WITH_AES_256_GCM_SHA384	CompressionMethods Length (01h)	NULL Compression (00h)
TPL	ExtensionsLength (MSB, 00h)	ExtensionsLength (LSB)	Extension elliptical_curves (MSB, 00h)	Extension elliptical_curves (LSB, 0Ah)
TPL	Elliptical_curves_length (MSB, 00h)	Elliptical_curves_length (LSB)	secp256r1 (NIST P-256) MSB (00h)	secp256r1 (NIST P-256) LSB. (17h)
TPL	brainpoolp256r1 MSB (00h)	brainpoolp256r1 LSB (1Ah)	brainpoolp384r1 MSB (00h)	brainpoolp384r1 LSB (1Bh)
TPL	Extension truncated_hmac (MSB, 00h)	Extension truncated_hmac (LSB, 04h)	ExtensionData_length (MSB 00h)	ExtensionData_len (LSB 00h))
TPL	Extension max.fragment_length (MSB 00h)	Extension max.fragment_length (LSB 01h)	Extension length (MSB 00h)	Extension Length (01h)
TPL	Extension Data (02h) Max.Fragment Size 1 kbyte	Extension encrypt_then_mac (MSB 00h)	Extension encrypt_then_mac (LSB 16h)	Extension length (MSB 00h)
TPL	Extension Length (LSB 00h)			

NOTE: This example contains 4 cipher suites to show possible applications but in practice at least one cipher suite is sufficient.

F.D.4 Second message flight – ServerHello, Certificate, ServerKeyExchange, CertificateRequest, ServerHelloDone (from server to client)

F.D.4.1 First TLS fragment

Table F.D.4 – Second message flight, First TLS fragment

Layer	0 (first byte)	1	2	3
DLL	L = nn	C = 53h/73h (SND-UD)	MFCT_0(GW)	MFCT_1(GW)
DLL	ID_0(GW)	ID_1(GW)	ID_2(GW)	ID_3(GW)
DLL	VER(GW)	DEVTYPE(31h)		
ELL	CI = 8Eh(ELL)	CC	ACC	MFCT_0 (OMS end-device)
ELL	MFCT_1(OMS end-device)	ID_0(OMS end-device)	ID_1(OMS end-device)	ID_2(OMS end-device)
ELL	ID_3(OMS end-device)	VER(OMS end-device)	DEVTYPE(OMS end-device)	
	CI = 90(AFL)	AFL.AFLL(05h)	AFL.FCL(01h) LSB	AFL.FCL(70h)MSB MF = 1, ML
AFL	AFL.MCL(40h) ML	AFL.ML (8Fh, LSB)	AFL.ML (01h, MSB)	
TPL	CI = 5Fh	ID_0(MTR)	ID_1(MTR)	ID_2(MTR)
TPL	ID_3(MTR)	MFCT_0(MTR)	MFCT_1(MTR)	VER(MTR)
TPL	DEVTYPE(MTR)	ACC	STS	CF[7..0]
TPL	CF[15..8](MODE 13)	CFE[7..0](TLSPROT)		
TPL	ContentType (16h, Handshake)	ProtoMajor (03h)	ProtoMinor (03h) TLS1.2	TLSSDULen (MSB)
TPL	TLSSDULen (LSB)	HSMmsgType (02h ServerHello)	HSMmsgLength (MSB, 00h)	HSMmsgLength (00h)
TPL	HSMmsgLength (LSB 26h)	HSServerVersionMajor (03h)	HSServerVersionMinor (03h) TLS1.2	HSServerRandom 32 (MSB)
TPL	HSServerRandom 31	HSServerRandom 30	HSServerRandom 29	HSServerRandom 28
TPL	HSServerRandom 27	HSServerRandom	HSServerRandom	HSServerRandom
TPL	HSServerRandom 23	HSServerRandom	HSServerRandom	HSServerRandom
TPL	HSServerRandom 19	HSServerRandom	HSServerRandom	HSServerRandom
TPL	HSServerRandom 15	HSServerRandom	HSServerRandom	HSServerRandom
TPL	HSServerRandom 11	HSServerRandom	HSServerRandom	HSServerRandom
TPL	HSServerRandom 7	HSServerRandom	HSServerRandom	HSServerRandom
TPL	HSServerRandom 3	HSServerRandom 2	HSServerRandom 1	HSSessionIDLength (00h) ²
TPL	CipherSuite (MSB, C0h) TLS_ECDHE_ECDSA_WITH_AES_128_CBC_SHA256	CipherSuite (LSB, 23h) TLS_ECDHE_ECDSA_WITH_AES_128_CBC_SHA256	CompressionSelected (NONE, 00h)	ExtensionsLength (MSB 00h)
TPL	ExtensionsLength (LSB)	Extension elliptical_curves (MSB, 00h)	Extension elliptical_curves (LSB, 0Ah)	Elliptical_curves_length (MSB, 00h)
TPL	Elliptical_curves_length (LSB 02)	secp256r1 (NIST P-256) MSB (00h)	secp256r1(NIST P-256) LSB. (17h)	Extension max.fragment_length

² **NOTE:** If TLS session resumption is used, this field contains the SessionID (up to 32 bytes)

Layer	0 (first byte)	1	2	3
				(MSB 00h)
TPL	Extension max.fragment_length (LSB 01h)	Extension length (MSB 00h)	Extension Length (01h)	Extension Data (02h) Max.Fragment Size 1 kbyte
TPL	Extension encrypt_then_mac (MSB 00h)	Extension encrypt_then_mac (LSB 16h)	Extension length (MSB 00h)	Extension Length (LSB 00h)
TPL	HSMMsgType (0Bh Certificate)	HSMMsgLength (MSB,00h)	HSMMsgLength (00h)	HSMMsgLength (LSB,F9h)
TPL	ASN.1 CertData (SEQUENCE 30h)	ASN.1 CertData (82h)	ASN.1 CertData	ASN.1 CertData
TPL
TPL	ASN.1 CertData	ASN.1 CertData	ASN.1 CertData	ASN.1 CertData

The OMS end-device answers with ACK.

F.D.4.2 Second TLS fragment

Table F.D.5 – Second message flight, second TLS fragment

Layer	0 (first byte)	1	2	3
DLL	L = nn	C = 53h/73h (SND-UD)	MFCT_0(GW)	MFCT_1(GW)
DLL	ID_0(GW)	ID_1(GW)	ID_2(GW)	ID_3(GW)
DLL	VER(GW)	DEVTYPE(31h)		
ELL	CI = 8Eh(ELL)	CC	ACC	MFCT_0 (OMS end-device)
ELL	MFCT_1(OMS end-device)	ID_0(OMS end-device)	ID_1(OMS end-device)	ID_2(OMS end-device)
ELL	ID_3(OMS end-device)	VER(OMS end-device)	DEVTYPE(OMS end-device)	
AFL	CI = 90h(AFL)	AFL.AFLL(02h)	AFL.FCL(02h) LSB	AFL.FCL(00h) MSB
TPL	ASN1.CertData	ASN1.CertData	ASN1.CertData	ASN1.CertData
TPL
TPL	ASN.1 CertData	ASN.1 CertData	ASN.1 CertData	ASN.1 CertData (CertDataEnd)
TPL	HSMmsgType (0Ch) ServerKeyExchange	HSMmsgLength (MSB, 00h)	HSMmsgLength (00h)	HSMmsgLength (LSB, 25h)
TPL	CurveType = namedCurve (03h)	00h (MSB)	17h (LSB, secp256r1)	Public ECC Key Point Len uncompress.(41h)
TPL	ECC Point format (04h)	ECC Point Public (2*32 bytes)	ECCPoint Public	ECCPoint Public
TPL	ECCPoint Public	ECCPoint Public	ECCPoint Public	ECCPoint Public
TPL
TPL	ECCPoint Public	HSMmsgType (0Dh) CertificateRequest	HSMmsgLength (MSB, 00h)	HSMmsgLength (00h)
TPL	HSMmsgLength (LSB 08h)	CertificateTypeList Len (01h)	40h ECDSA_Sign	00h (MSB Signature HashAlgoListLen)
TPL	02h (LSB Signature HashAlgoListLen)	04h (SHA256)	04h (SHA256) 03h (ECDSA)	CAListLen (MSB 00h)
TPL	CAListLen (LSB 00h)	HSMmsgType (0Eh) ServerHelloDone	HSMmsgLength (MSB, 00h)	HSMmsgLength (00h)
TPL	HSMmsgLength (LSB, 00h)			

The OMS end-device answers with ACK. The OMS end-device waits for REQ-UD2 from the gateway.

F.D.5 Third message flight, Certificate, ClientKeyExchange, CertificateVerify, ChangeCipherSpec, Finished (from client to server)

F.D.5.1 First TLS fragment

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Table F.D. 6 – Third message flight, first TLS fragment

Layer	0 (first byte)	1	2	3
DLL	L = nnh	C = 08h/28h (RSP-UD)	MFCT_0(MTR)	MFCT_1(MTR)
DLL	ID_0(MTR)	ID_1(MTR)	ID_2(MTR)	ID_3(MTR)
DLL	VER(MTR)	DEVTYPE(MTR)		
ELL	CI = 8Eh(ELL)	CC	ACC	MFCT_0 (GW)
ELL	MFCT_1(GW)	ID_0(GW)	ID_1(GW)	ID_2(GW)
ELL	ID_3(GW)	VER(GW)	DEVTYPE(GW)	
AFL	CI = 90h(AFL)	AFL.AFLL(05h)	AFL.FCL(01h) LSB	AFL.FCL(70h) MSB MF = 1, MCL, ML
AFL	AFL.MCL(40h), ML	AFL.ML (B8h, LSB)	AFL.ML (01h, MSB)	
TPL	CI = 9Eh	ACC	STS	CF[7..0]
TPL	CF[15..8] (MODE 13)	CFE[7..0] (TLSPROT)		
TPL	ContentType (16h, Handshake)	ProtoMajor (03h)	ProtoMinor (03h) TLS1.2	TLSSDULen (MSB)
TPL	TLSSDULen (LSB)	HSMType (0Bh) Certificate	HSMLength (MSB)	HSMLength
TPL	HSMLength (LSB)	ASN.1 CertData (SEQUENCE 30h)	ASN.1 CertData (82h)	ASN.1 CertData
TPL	ASN1.CertData	ASN1.CertData	ASN1.CertData	ASN1.CertData
TPL
TPL	ASN.1 CertData	ASN.1 CertData	ASN.1 CertData	ASN.1 CertData

The OMS end-device waits for REQ-UD2 from the gateway.

Layer	0 (first byte)	1	2	3
TPL	ProtoMinor (03h) TLS1.2	TLSSDULen (MSB, 00h)	TLSSDULen (LSB, 34h)	HSMType (14h) Finished
TPL	HSMLength (MSB, 00h)	HSMLength	HSMLength (LSB, 30h)	Encrypted HSHash
TPL	Encrypted HSHash	Encrypted HSHash	Encrypted HSHash	Encrypted HSHash
TPL	Encrypted HSHash	Encrypted HSHash	Encrypted HSHash	Encrypted HSHash
TPL	Encrypted HSHash	Encrypted HSHash	Encrypted HSHash	Encrypted Padding(03h)
TPL	Encrypted Padding(03h)	Encrypted Padding(03h)	Encrypted PadLength(03h)	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	

The OMS end-device waits for SND-UD from the gateway.

NOTE 1: The handshake hash is not truncated, if MAC truncation is negotiated.

NOTE 2: Due to the length of the certificate it might be necessary to use a third fragment.

F.D.6 Fourth message flight – ChangeCipherSpec, finished (from server to client)

Using cipher suite AES128-CBC SHA256 with encrypt-then-mac, without truncated HMAC extensions.

5

Table F.D.8 – Fourth message flight

Layer	0 (first byte)	1	2	3
DLL	L = nn	C = 53h/73h (SND-UD)	MFCT_0(GW)	MFCT_1(GW)
DLL	ID_0(GW)	ID_1(GW)	ID_2(GW)	ID_3(GW)
DLL	VER(GW)	DEVTYPE(31h)		
ELL	CI = 8Eh(ELL)	CC	ACC	MFCT_0(OMS end-device)
ELL	MFCT_1(OMS end-device)	ID_0(OMS end-device)	ID_1(OMS end-device)	ID_2 (OMS end-device)
ELL	ID_3 (OMS end-device)	VER(OMS end-device)	DEVTYPE(OMS end-device)	
TPL	CI = 5Fh	ID_0(MTR)	ID_1(MTR)	ID_2(MTR)
TPL	ID_3(MTR)	MFCT_0(MTR)	MFCT_1(MTR)	VER(MTR)
TPL	DEVTYPE(MTR)	ACC	STS	CF[7..0]
TPL	CF[15..8](MODE 13,)	CF [23..16](TLSHS)		
TPL	ContentType (14h, ChangeCipherSpec)	ProtoMajor (03h)	ProtoMinor (03h) TLS1.2	TLSSDULen (MSB, 00h)
TPL	TLSSDULen (LSB, 05h)	HSMmsgType (01h) ChangeCipherSpec	HSMmsgLength (MSB, 00h)	HSMmsgLength
TPL	HSMmsgLength (LSB, 01h)	Fix parameter (01h)	ContentType (16h, Handshake)	ProtoMajor (03h)
TPL	ProtoMinor (03h) TLS1.2	TLSSDULen (MSB, 00h)	TLSSDULen (LSB, 34h)	HSMmsgType (14h) Finished
TPL	HSMmsgLength (MSB, 00h)	HSMmsgLength	HSMmsgLength (LSB, 30h)	Encrypted HSHash
TPL	Encrypted HSHash	Encrypted HSHash	Encrypted HSHash	Encrypted HSHash
TPL	Encrypted HSHash	Encrypted HSHash	Encrypted HSHash	Encrypted HSHash
TPL	Encrypted HSHash	Encrypted HSHash	Encrypted HSHash	Encrypted Padding(03h)
TPL	Encrypted Padding(03h)	Encrypted Padding(03h)	Encrypted PadLength(03h)	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	

The OMS end-device answers with ACK.

F.D.7 Xth message flight on wireless M-Bus – M-Bus application data transfer (from server to client)

This is an example of an M-Bus application command protected by AES128-CBC SHA256 transferred via the wireless M-Bus.

Table F.D.9 – Application data (wireless M-Bus, SND-UD)

Layer	0 (first byte)	1	2	3
DLL	L = nn	C = 53h/73h (SND-UD)	MFCT_0(GW)	MFCT_1(GW)
DLL	ID_0(GW)	ID_1(GW)	ID_2(GW)	ID_3(GW)
DLL	VER(GW)	DEVTYPE(31h)		
ELL	CI = 8Ch(ELL)	CC	ACC	
TPL	CI = 5Bh	ID_0(MTR)	ID_1(MTR)	ID_2(MTR)
TPL	ID_3(MTR)	MFCT_0(MTR)	MFCT_1(MTR)	VER(MTR)
TPL	DEVTYPE(MTR)	ACC	STS	CF[7..0]
TPL	CF[15..8](MODE 13)	CFE[7..0](TLSPROT)		
TPL	ContentType (17h, Application Data)	ProtoMajor (03h)	ProtoMinor (03h) TLS1.2	TLSSDULen (MSB, 00h)
TPL	TLSSDULen (LSB, 40h) incl. IV, HMAC and padding	IV (random)	IV (random)	IV (random)
TPL	IV (random)	IV (random)	IV (random)	IV (random)
TPL	IV (random)	IV (random)	IV (random)	IV (random)
TPL	IV (random)	IV (random)	IV (random)	IV (random)
APL	Encrypted AppData DIF	Encrypted AppData VIF	Encrypted AppData Value	Encrypted AppData Value
APL	Encrypted AppData Value	Encrypted AppData Value	Encrypted AppData DIF2	Encrypted AppData VIF2
APL	Encrypted AppData VIFE	Encrypted AppData Value	Encrypted AppData Value	Encrypted AppData Value
APL	Encrypted AppData Value	Encrypted Padding(02h)	Encrypted Padding(02h)	Encrypted PadLength(02h)
TPL	SHA256 HMAC (first byte; shown without truncation)	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC 32 (last byte)

The OMS end-device answers with ACK. The gateway has to send a REQ-UD2 to receive application data from the OMS end-device.

The HMAC size for SHA256 is 32 bytes. If both sides (server and client) support HMAC truncation (see RFC6066), the HMAC size can be truncated to 10 bytes.

The padding value and PadLen is calculated according to [RFC5246] 3.2.3.2.

Due to the 16 byte block length requirement for AES128 a minimum of 0-15 padding bytes is added before the PadLen byte. To hide the actual payload size, a multiple of 16 bytes padding can be additionally added (up to a total of 255 padding bytes).

Example: AppData Length = 13, HMAC = 32, PadLen = 1, IVLen = 16: $(13+32+16+1) \text{ MOD } 16 = 14$. Thus the padding value and PadLen can have the values $(16-14) = 0x02, 0x12, 0x22, 0x32 \dots 0xF2$.

Because the link layer (12 byte), extended link layer (3 or 11 byte), the optional AFL header (7 byte) and the TPL header (0,5 or 13 byte) adds 19 to 43 bytes, the APL data size is typically around 200 bytes. It is the task of the application layer to discover the maximum unfragmented TPL+APL size.

NOTE: Because of the HMAC, padding and IV overhead it is more efficient to do the fragmentation in the AFL layer than within the TLS Application Layer.

F.D.8 Xth message flight on wired M-Bus – M-Bus application data transfer (from client to server)

This is an example of an M-Bus application data protected by AES128-CBC SHA256 transferred via wired M-Bus.

- 5 The server sends a request of data.

Table F.D.10 – Application data (wired M-Bus, REQ-UD2)

Layer	0 (first byte)	1	2	3
DLL	10h	C = 5B/7Bh (REQ-UD)	PAddr (Primary M-Bus addr.)	Chksum
DLL	0x16			

The client responds the currently selected application data.

Table F.D.11 – Application data (wired M-Bus, RSP-UD)

Layer	0 (first byte)	1	2	3
DLL	68h	L = nn	L = nn	68h
DLL	C = 08h (RSP-UD)	PAddr (Primary M-Bus addr.)		
TPL	CI = 72h	ID_0(MTR)	ID_1(MTR)	ID_2(MTR)
TPL	ID_3(MTR)	MFCT_0(MTR)	MFCT_1(MTR)	VER(MTR)
TPL	DEVTYPE(MTR)	ACC	STS	CF[7..0]
TPL	CF[15..8](MODE 13)	CFE[7..0](TLSPROT)		
TPL	ContentType (17h, Application Data)	ProtoMajor (03h)	ProtoMinor (03h) TLS1.2	TLSSDULen (MSB, 00h)
TPL	TLSSDULen (LSB, 40h) incl. IV, HMAC and padding			
TPL	IV (random)	IV (random)	IV (random)	IV (random)
TPL	IV (random)	IV (random)	IV (random)	IV (random)
TPL	IV (random)	IV (random)	IV (random)	IV (random)
TPL	IV (random)	IV (random)	IV (random)	IV (random)
TPL	Encrypted AppData DIF	Encrypted AppData VIF	Encrypted AppData Value	Encrypted AppData Value
APL	Encrypted AppData Value	Encrypted AppData Value	Encrypted AppData DIF2	Encrypted AppData VIF2
APL	Encrypted AppData VIFE	Encrypted AppData Value	Encrypted AppData Value	Encrypted AppData Value
APL	Encrypted AppData Value	Encrypted Padding(02h)	Encrypted Padding(02h)	Encrypted PadLength(02h)
TPL	SHA256 HMAC (first byte; shown without truncation)	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC 32(LastByte)
DLL	Chksum	0x16		

Appendix F.E (informative): Examples of the Security Information Transfer Protocol

This appendix shows example datagrams between a gateway and an OMS end-device transporting the Security Information Transport Protocol (SITP).

F.E.1 Example Master key renewal with security mode 13 (bidirectional communication) – key transfer

The following table shows a SND-UD datagram sent by the gateway after successful reception of a SND-NR datagram by the OMS end-device. Thereby, it starts a bidirectional communication between both devices. Goal of the bidirectional communication is to transport and activate a new AES128 encryption key (actually the master key) to the OMS end-device.

Table F.E.1 – SITP transfer key command

Layer	0 (first byte)	1	2	3
DLL	L = nn	C = 43h (SND-UD2)	MFCT_0(GW)	MFCT_1(GW)
DLL	ID_0(GW)	ID_1(GW)	ID_2(GW)	ID_3(GW)
DLL	VER(GW)	DEVTYPE(31h)		
ELL	CI = 8Ch(ELL)	CC	ACC	
TPL	CI = C3h	ID_0(MTR)	ID_1(MTR)	ID_2(MTR)
TPL	ID_3(MTR)	MFCT_0(MTR)	MFCT_1(MTR)	VER(MTR)
TPL	DEVTYPE(MTR)	ACC	STS	CF[7..0]
TPL	CF[15..8](MODE 13)	CFE[7..0](TLSPROT)		
TPL	ContentType(17h, Application Data)	ProtoMajor (03h)	ProtoMinor(03h) TLS1.2	TLSSDULen (MSB, 00h)
TPL	TLSSDULen (LSB, 60h) incl. IV, HMAC and padding			
TPL	IV (random)	IV (random)	IV (random)	IV (random)
TPL	IV (random)	IV (random)	IV (random)	IV (random)
TPL	IV (random)	IV (random)	IV (random)	IV (random)
TPL	IV (random)	IV (random)	IV (random)	IV (random)
APL-SITP	Block length (BL) LSB = 26h (38 bytes)	Block length (BL) MSB = 00h	Block ID (BID) = 00h	Block Control Field (BCF) = 00h (Transfer security information)
APL-SITP	Recipient ID = 00h (no dedicated application)	Data Structure Identifier (DSI) = 01h (Wrapping AES128 according NIST SP 800-38F type KWP)	Data structure header (DSH1) = FFh (WrapperKey ID = FFh identifies no wrapping in the TLS channel)	Data structure header (DSH2) = FFh (WrapperKeyVersion = FFh identifies no WrapperKey)
APL-SITP	Wrapped struct. 01 (ICV) = A6h	Wrapped struct. 01 (ICV) = 59h	Wrapped struct. 01 (ICV) = 59h	Wrapped struct. 01 (ICV) = A6h
APL-SITP	Wrapped struct. 01 (MLI) = 00h (MSB)	Wrapped struct. 01 (MLI) = 00h	Wrapped struct. 01 (MLI) = 00h	Wrapped struct. 01 (MLI) = 17h (LSB)
APL-SITP	Key/random z1 (MSB)	Key/random z1	Key/random z1	Key/random z1
APL-SITP	Key/random z1	Key/random z1	Key/random z1	Key/random z1

Layer	0 (first byte)	1	2	3
APL-SITP	Key/random z1	Key/random z1	Key/random z1	Key/random z1
APL-SITP	Key/random z1	Key/random z1	Key/random z1	Key/random z1 (LSB)
APL-SITP	Target Time = 00h (LSB)	Target Time = 00h	Target Time = 00h	Target Time = 80h
APL-SITP	Target Time = 30h (MSB)	KeyID = 00h (master key)	KeyVersion = 01h (First replaced key)	Padding = 00h
TPL	Encrypted Padding(07h)	Encrypted Padding(07h)	Encrypted Padding(07h)	Encrypted Padding(07h)
TPL	Encrypted Padding(07h)	Encrypted Padding(07h)	Encrypted Padding(07h)	Encrypted PadLength(07h)
TPL	SHA256 HMAC (first byte; shown without truncation)	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC 32 (last byte)

Assuming the OMS end-device receives the datagrams and processes them successfully, the corresponding response by the OMS end-device is as follows:

Table F.E.2 – SITP transfer key command response

Layer	0 (first byte)	1	2	3
DLL	L = nn	C = 08h (RSP-UD)	MFCT_0(MTR)	MFCT_1(MTR)
DLL	ID_0(MTR)	ID_1(MTR)	ID_2(MTR)	ID_3(MTR)
DLL	VER(MTR)	DEVTYPE(MTR)		
ELL	CI = 8Ch(ELL)	CC	ACC	
TPL	CI = C4h	ACC	STS	CF[7..0]
TPL	CF[15..8](MODE 13)	CFE[7..0](TLSPROT)		
TPL	ContentType(17h, Application Data)	ProtoMajor (03h)	ProtoMinor(03h) TLS1.2	TLSSDULen (MSB, 00h)
TPL	TLSSDULen (LSB, 40h) incl. IV, HMAC and padding			
TPL	IV (random)	IV (random)	IV (random)	IV (random)
TPL	IV (random)	IV (random)	IV (random)	IV (random)
TPL	IV (random)	IV (random)	IV (random)	IV (random)
TPL	IV (random)	IV (random)	IV (random)	IV (random)
APL-SITP	Block length (BL) LSB = 07h (7 bytes)	Block length (BL) MSB = 00h	Block ID (BID) = 00h	Block Control Field (BCF) = 80h (Status Response to Transfer security information)
APL-SITP	Recipient ID = 00h (no dedicated application)	Data Structure Identifier (DSI) = 22h (Status Response Structure)	Data structure header (DSH1) = FFh (KeyID = FFh identifies no wrapping in the TLS channel)	Data structure header (DSH2) = FFh (KeyVersion = FFh identifies no wrapper key)
APL-SITP	Block parameters = 00h (Successful command)	Encrypted Padding(06h)	Encrypted Padding(06h)	Encrypted Padding(06h)
TPL	Encrypted Padding(06h)	Encrypted Padding(06h)	Encrypted Padding(06h)	Encrypted PadLength(06h)
TPL	SHA256 HMAC (first byte; shown without truncation)	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC 32 (last byte)

F.E.2 Example Master key renewal with security mode 13 (bidirectional communication) – key activation/deactivation

At some point in time, the gateway activates the already transferred key by using the “combined activate/deactivate key” command. The corresponding datagram is as follows:

Table F.E.3 – SITP combined activate/deactivate key command

Layer	0 (first byte)	1	2	3
DLL	L = nn	C = 43h (SND-UD2)	MFCT_0(GW)	MFCT_1(GW)
DLL	ID_0(GW)	ID_1(GW)	ID_2(GW)	ID_3(GW)
DLL	VER(GW)	DEVTYPE(31h)		
ELL	CI = 8Ch(ELL)	CC	ACC	
TPL	CI = C3h	ID_0(MTR)	ID_1(MTR)	ID_2(MTR)
TPL	ID_3(MTR)	MFCT_0(MTR)	MFCT_1(MTR)	VER(MTR)
TPL	DEVTYPE(MTR)	ACC	STS	CF[7..0]
TPL	CF[15..8](MODE 13)	CFE[7..0](TLSPROT)		
TPL	ContentType(17h, Application Data)	ProtoMajor (03h)	ProtoMinor(03h) TLS1.2	TLSSDULen (MSB, 00h)
TPL	TLSSDULen (LSB, 60h) incl. IV, HMAC and padding			
TPL	IV (random)	IV (random)	IV (random)	IV (random)
TPL	IV (random)	IV (random)	IV (random)	IV (random)
TPL	IV (random)	IV (random)	IV (random)	IV (random)
TPL	IV (random)	IV (random)	IV (random)	IV (random)
APL-SITP	Block length (BL) LSB = 1Eh (30 bytes)	Block length (BL) MSB = 00h	Block ID (BID) = 00h	Block Control Field (BCF) = 04h (Combined activation/deactivation security information)
APL-SITP	Recipient ID = 00h (no dedicated application)	Data Structure Identifier (DSI) = 03h (Wrapping AES128 according NIST SP 800-38F type KWP)	Data structure header (DSH1) = FFh (WrapperKey ID = FFh identifies no wrapping in the TLS channel)	Data structure header (DSH2) = FFh (WrapperKeyVersion = FFh identifies no wrapper key)
APL-SITP	Wrapped struct. 03 (ICV) = A6h (MSB)	Wrapped struct. 03 (ICV) = 59h	Wrapped struct. 03 (ICV) = 59h	Wrapped struct. 03 (ICV) = A6h (LSB)
APL-SITP	Wrapped struct. 03 (MLI) = 00h (MSB)	Wrapped struct. 03 (MLI) = 00h	Wrapped struct. 03 (MLI) = 00h	Wrapped struct. 03 (MLI) = 0Ah (LSB)
APL-SITP	Target Time = 00h (LSB)	Target Time = 00h	Target Time = 00h	Target Time = 00h
APL-SITP	Target Time = 30h (MSB)	Activated KeyID = 00h (master key)	Activated KeyVersion = 01h (Activate previously transferred new master key version)	Deactivated KeyID = 00h (master key)
APL-SITP	Deactivated KeyVersion = 00h (Deactivate old master key version)	Option = 01h (no MessageCounter reset)	Padding = 00h	Padding = 00h
APL-SITP	Padding = 00h	Padding = 00h	Padding = 00h	Padding = 00h
TPL	Encrypted Padding(0Fh)	Encrypted Padding(0Fh)	Encrypted Padding(0Fh)	Encrypted Padding(0Fh)
TPL	Encrypted Padding(0Fh)	Encrypted Padding(0Fh)	Encrypted Padding(0Fh)	Encrypted Padding(0Fh)
TPL	Encrypted Padding(0Fh)	Encrypted Padding(0Fh)	Encrypted Padding(0Fh)	Encrypted Padding(0Fh)
TPL	Encrypted Padding(0Fh)	Encrypted Padding(0Fh)	Encrypted Padding(0Fh)	Encrypted PadLength(0Fh)

Layer	0 (first byte)	1	2	3
<i>TPL</i>	SHA256-HMAC (first byte; shown without truncation)	HMAC	HMAC	HMAC
<i>TPL</i>	HMAC	HMAC	HMAC	HMAC
<i>TPL</i>	HMAC	HMAC	HMAC	HMAC
<i>TPL</i>	HMAC	HMAC	HMAC	HMAC
<i>TPL</i>	HMAC	HMAC	HMAC	HMAC
<i>TPL</i>	HMAC	HMAC	HMAC	HMAC
<i>TPL</i>	HMAC	HMAC	HMAC	HMAC
<i>TPL</i>	HMAC	HMAC	HMAC	HMAC 32 (last byte)

The corresponding response of the OMS end-device to the “combined activate/deactivate key” command is as follows, assuming that the command is processed successfully.

Table F.E.4 – SITP combined activate/deactivate key command response

Layer	0 (first byte)	1	2	3
DLL	L = nn	C = 08h (RSP-UD)	MFCT_0(MTR)	MFCT_1(MTR)
DLL	ID_0(MTR)	ID_1(MTR)	ID_2(MTR)	ID_3(MTR)
DLL	VER(MTR)	DEVTYPE(MTR)		
ELL	CI = 8Ch(ELL)	CC	ACC	
TPL	CI = C4h	ACC	STS	CF[7..0]
TPL	CF[15..8](MODE 13)	CFE[7..0](TLSPROT)		
TPL	ContentType(17h, Application Data)	ProtoMajor (03h)	ProtoMinor(03h) TLS1.2	TLSSDULen (MSB, 00h)
TPL	TLSSDULen (LSB, 40h) incl. IV, HMAC and padding			
TPL	IV (random)	IV (random)	IV (random)	IV (random)
TPL	IV (random)	IV (random)	IV (random)	IV (random)
TPL	IV (random)	IV (random)	IV (random)	IV (random)
TPL	IV (random)	IV (random)	IV (random)	IV (random)
APL-SITP	Block length (BL) LSB = 07h	Block length (BL) MSB = 00h	Block ID (BID) = 00h	Block Control Field (BCF) = 84h (Status Response to Combined activation/deactivation security information)
APL-SITP	Recipient ID = 00h (no dedicated application)	Data Structure Identifier (DSI) = 22h (Status Response Structure)	Data structure header (DSH1) = FFh (KeyID = FFh identifies no wrapping in the TLS channel)	Data structure header (DSH2) = FFh (KeyVersion = FFh identifies no wrapper key)
APL-SITP	Block parameters = 00h (Successful command)	Encrypted Padding(06h)	Encrypted Padding(06h)	Encrypted Padding(06h)
TPL	Encrypted Padding(06h)	Encrypted Padding(06h)	Encrypted Padding(06h)	Encrypted PadLength(06h)
TPL	SHA256-HMAC (first byte; shown without truncation)	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC
TPL	HMAC	HMAC	HMAC	HMAC 32 (last byte)

Appendix F.F (informative): Example certificate

An example of a self-signed certificate for a gas meter is shown below.

NOTE 1: The size of the certificate below is 339 bytes.

NOTE 2: An X509v3 Key Usage “Digital Signature” is sufficient for use with ECDHE cipher suites, because TLS ephemeral encryption key pairs are generated independent from the certificate keys.

Certificate Data (TEXT with data fields):

Version: 3 (0x2)

Serial Number: c9:e7:c5:e1:44:ab:20:ba

Signature Algorithm: ecdsa-with-SHA256

Issuer: CN = 7mfc0100001234.mtr

Validity

Not Before: Aug 1 00:00:00 2019 GMT

Not After : Aug 1 00:00:00 2026 GMT

Subject: CN = 7mfc0100001234.mtr

Subject Public Key Info:

Public Key Algorithm: id-ecPublicKey

Public-Key: (256 bit)

pub:

04:8a:43:29:95:15:fa:fb:9f:45:6a:73:7d:b7:05:be:ee:15:
d7:1e:cc:68:22:aa:c2:f4:ca:12:71:e4:58:8e:cc:76:35:5b:
b4:90:06:35:d2:ee:56:29:49:12:72:06:23:75:f7:90:79:f6:
48:74:07:92:dc:96:f0:ac:92:90:d4

ASN1 OID: brainpoolP256r1

X509v3 extensions:

X509v3 Basic Constraints:

CA:TRUE

X509v3 Key Usage: Digital Signature

Signature Algorithm: ecdsa-with-SHA256

30:46:02:21:00:82:83:d2:87:08:69:a1:c8:20:96:8a:96:66:3d:c9:15:9b:81:82:9f:
96:d5:fc:41:e1:27:22:cc:71:4d:e1:97:02:21:00:88:99:82:c0:17:03:15:1c:a2:11:
1c:9b:34:a6:0b:9d:d2:a6:0f:2b:32:79:b3:57:41:9e:5a:e9:1a:be:79:e2

Certificate Data (DER):

```
0000000: 30 82 01 4f 30 81 f5 a0 03 02 01 02 02 09 00 c9
0000010: e7 c5 e1 44 ab 20 ba 30 0a 06 08 2a 86 48 ce 3d
0000020: 04 03 02 30 1d 31 1b 30 19 06 03 55 04 03 0c 12
5 0000030: 37 6f 6d 73 31 61 30 30 30 30 30 34 30 30 2e 6d
0000040: 74 72 30 1e 17 0d 31 39 30 36 30 36 30 37 32 36
0000050: 32 32 5a 17 0d 32 36 30 36 30 34 30 37 32 36 32
0000060: 32 5a 30 1d 31 1b 30 19 06 03 55 04 03 0c 12 37
0000070: 6f 6d 73 31 61 30 30 30 30 30 34 30 30 2e 6d 74
10 0000080: 72 30 5a 30 14 06 07 2a 86 48 ce 3d 02 01 06 09
0000090: 2b 24 03 03 02 08 01 01 07 03 42 00 04 8a 43 29
00000a0: 95 15 fa fb 9f 45 6a 73 7d b7 05 be ee 15 d7 1e
00000b0: cc 68 22 aa c2 f4 ca 12 71 e4 58 8e cc 76 35 5b
00000c0: b4 90 06 35 d2 ee 56 29 49 12 72 06 23 75 f7 90
15 00000d0: 79 f6 48 74 07 92 dc 96 f0 ac 92 90 d4 a3 1d 30
00000e0: 1b 30 0c 06 03 55 1d 13 04 05 30 03 01 01 ff 30
00000f0: 0b 06 03 55 1d 0f 04 04 03 02 07 80 30 0a 06 08
0000100: 2a 86 48 ce 3d 04 03 02 03 49 00 30 46 02 21 00
0000110: 82 83 d2 87 08 69 a1 c8 20 96 8a 96 66 3d c9 15
20 0000120: 9b 81 82 9f 96 d5 fc 41 e1 27 22 cc 71 4d e1 97
0000130: 02 21 00 88 99 82 c0 17 03 15 1c a2 11 1c 9b 34
0000140: a6 0b 9d d2 a6 0f 2b 32 79 b3 57 41 9e 5a e9 1a
0000150: be 79 e2
```